

# *Quantum Circuits*

School on Quantum Computing @Yagami

Day 1, Lesson 5

16:00-17:00, March 22, 2005

Eisuke Abe

Department of Applied Physics and Physico-Informatics,  
and CREST-JST, Keio University



# Outline

- Bloch sphere representation
- Rotation gates
- Universality proof
  - An arbitrary controlled- $U$  gate can be implemented using only single qubit gates and CNOT
  - An arbitrary (controlled) $^n$ - $U$  gate can be implemented using single qubit gates and CNOT
  - Two-level unitary gates are universal
  - Single-qubit gates and CNOT are universal
  - Hadamard,  $S$ ,  $T$ , and CNOT are universal

# Bloch sphere representation

$$|\psi\rangle = \alpha|0\rangle + \beta|1\rangle$$

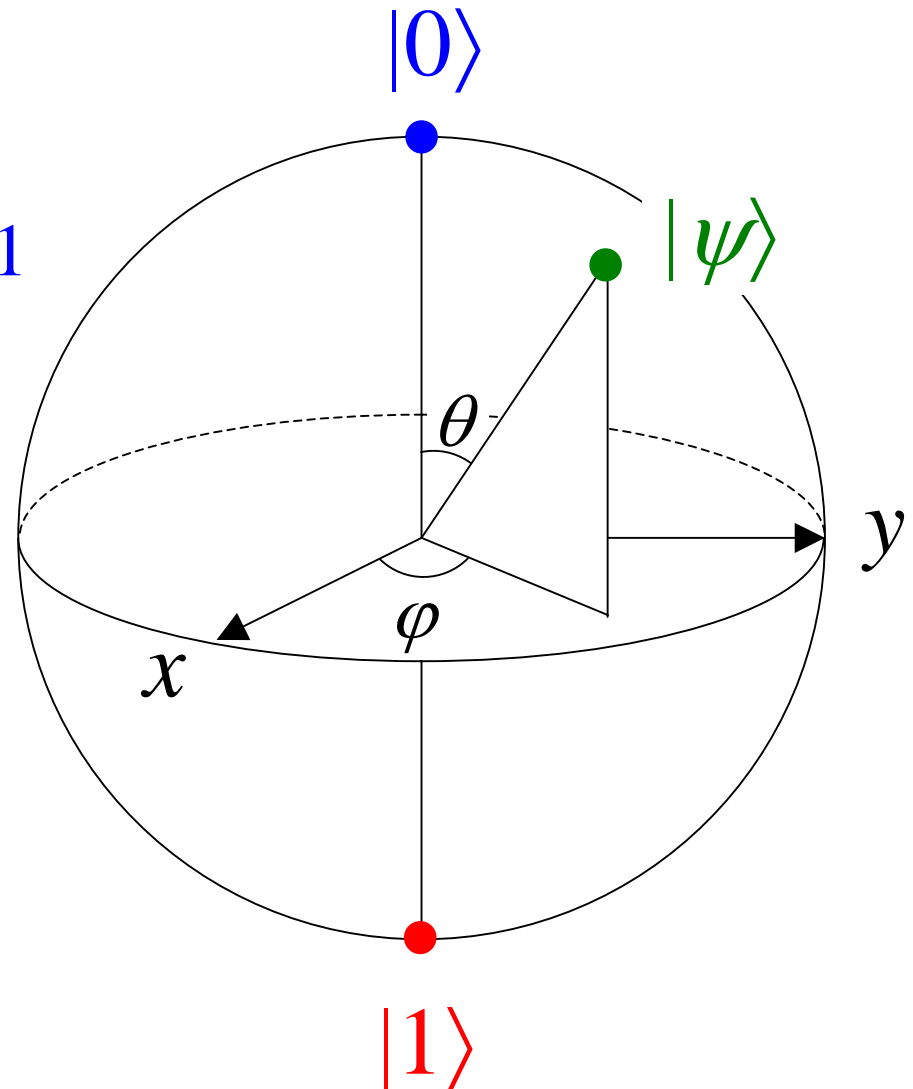
↓  $|\alpha|^2 + |\beta|^2 = 1$

$$|\psi\rangle = e^{i\gamma} \left( \cos \frac{\theta}{2} |0\rangle + e^{i\varphi} \sin \frac{\theta}{2} |1\rangle \right)$$

No  
observable  
effect

↓

$$|\psi\rangle = \cos \frac{\theta}{2} |0\rangle + e^{i\varphi} \sin \frac{\theta}{2} |1\rangle$$



# Important single qubit gates

$$X = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix} \quad Y = \begin{bmatrix} 0 & -i \\ i & 0 \end{bmatrix} \quad Z = \begin{bmatrix} 1 & 0 \\ 0 & -1 \end{bmatrix}$$

$$H = \frac{1}{\sqrt{2}} \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix} \quad S = \begin{bmatrix} 1 & 0 \\ 0 & i \end{bmatrix} \quad T = \begin{bmatrix} 1 & 0 \\ 0 & e^{i\pi/4} \end{bmatrix}$$

$$X^2 = Y^2 = Z^2 = H^2 = I$$

$$S = T^2, S^2 = Z$$

$$[X, Y] = 2iZ, \{X, Y\} = 0, \dots$$

$$HXH = Z, HYH = -Y, HZH = X$$

# Exponential operator

$$\exp(iAx) \equiv \sum_{n=0}^{\infty} \frac{(iAx)^n}{n!}$$

$$A^2 = I \quad \Rightarrow \quad \exp(iAx) = \cos x \cdot I + i \sin x \cdot A$$

This operator is important because it appeared in the solution to Schrödinger equation

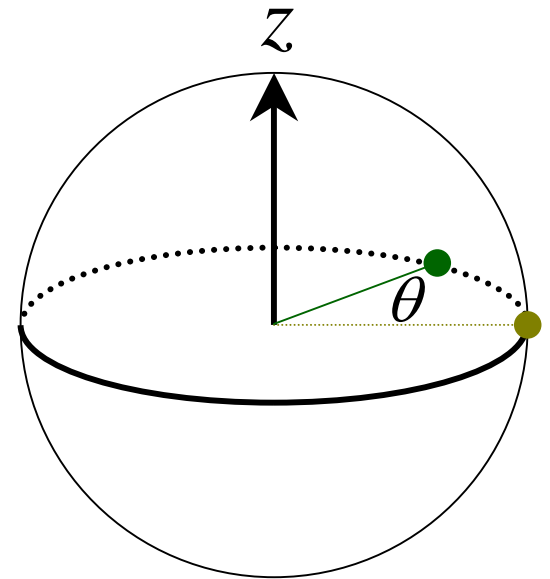
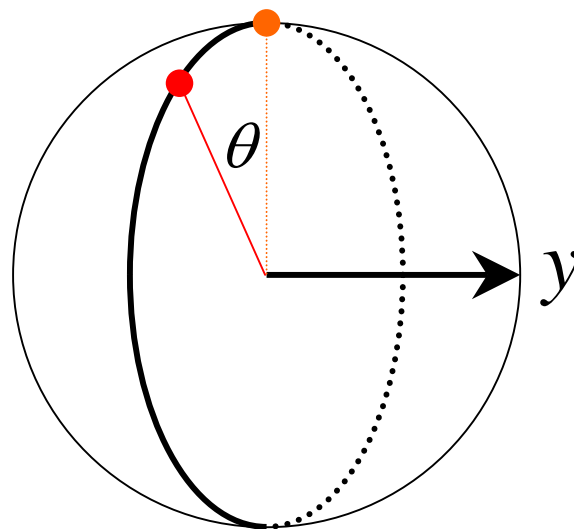
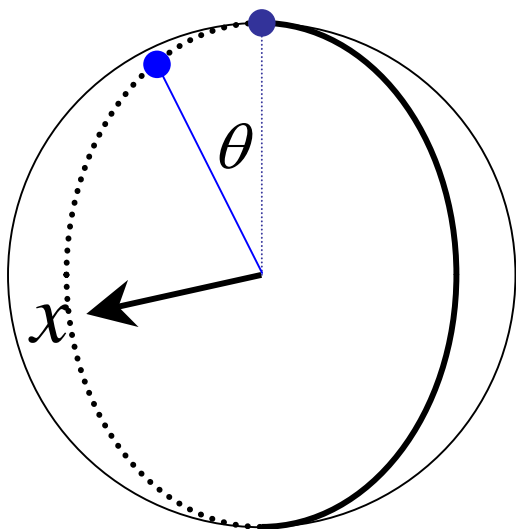
$$|\psi(t + \Delta t)\rangle = \exp\left[\frac{-iH\Delta t}{\hbar}\right] |\psi(t)\rangle$$

# Rotation gates

$$R_x(\theta) \equiv e^{-i\theta X/2} = \cos\frac{\theta}{2}I - i\sin\frac{\theta}{2}X = \begin{bmatrix} \cos(\theta/2) & -i\sin(\theta/2) \\ -i\sin(\theta/2) & \cos(\theta/2) \end{bmatrix}$$

$$R_y(\theta) \equiv e^{-i\theta Y/2} = \cos\frac{\theta}{2}I - i\sin\frac{\theta}{2}Y = \begin{bmatrix} \cos(\theta/2) & -\sin(\theta/2) \\ \sin(\theta/2) & \cos(\theta/2) \end{bmatrix}$$

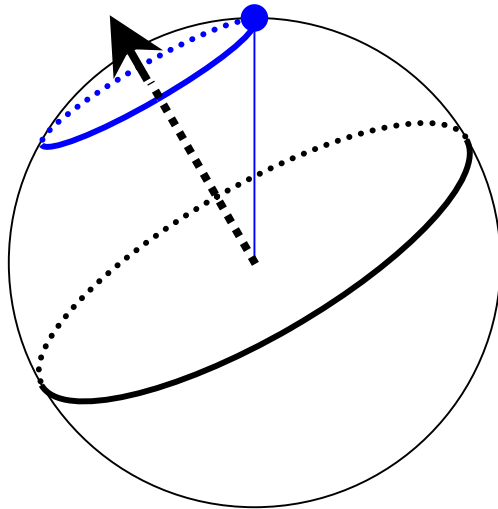
$$R_z(\theta) \equiv e^{-i\theta Z/2} = \cos\frac{\theta}{2}I - i\sin\frac{\theta}{2}Z = \begin{bmatrix} \exp(-i\theta/2) & 0 \\ 0 & \exp(i\theta/2) \end{bmatrix}$$



# Rotation about the $\hat{n}$ axis

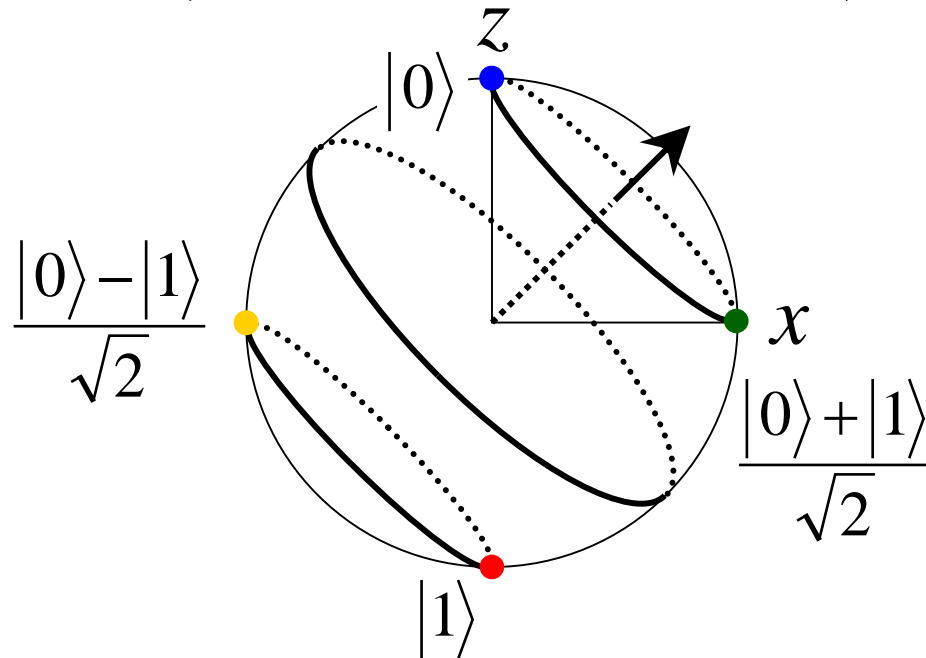
$$R_{\hat{n}}(\theta) \equiv \exp(-i\theta \hat{n} \cdot \hat{\sigma} / 2) = \cos \frac{\theta}{2} I - i \sin \frac{\theta}{2} (n_x X + n_y Y + n_z Z)$$

$$\hat{n} = \begin{pmatrix} n_x \\ n_y \\ n_z \end{pmatrix} \quad \hat{\sigma} = \begin{pmatrix} X \\ Y \\ Z \end{pmatrix}$$



## Example

$$H = \frac{X + Z}{\sqrt{2}} \Rightarrow \theta = \pi, \quad \hat{n} = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 \\ 0 \\ 1 \end{pmatrix}$$



# Road to universality proof

1. An arbitrary controlled- $U$  gate can be implemented using only single qubit gates and CNOT
2. An arbitrary (controlled) $^n$ - $U$  gate can be implemented using single qubit gates and CNOT
3. Two-level unitary gates are universal
4. Single-qubit gates and CNOT are universal
5. Hadamard,  $S$ ,  $T$ , and CNOT are universal



# Road to universality proof

1. An arbitrary controlled- $U$  gate can be implemented using only single qubit gates and CNOT
2. An arbitrary (controlled) $^n$ - $U$  gate can be implemented using single qubit gates and CNOT
3. Two-level unitary gates are universal
4. Single-qubit gates and CNOT are universal
5. Hadamard,  $S$ ,  $T$ , and CNOT are universal

# Z-Y decomposition

For an arbitrary single qubit gate  $U$ , there exist real numbers  $\alpha$ ,  $\beta$ ,  $\gamma$ , and  $\delta$  such that

$$U = e^{i\alpha} R_z(\beta) R_y(\gamma) R_z(\delta)$$

Proof

$$\begin{aligned} U &= \begin{bmatrix} e^{i(\alpha-\beta/2-\delta/2)} \cos \frac{\gamma}{2} & -e^{i(\alpha-\beta/2+\delta/2)} \sin \frac{\gamma}{2} \\ e^{i(\alpha+\beta/2-\delta/2)} \sin \frac{\gamma}{2} & e^{i(\alpha+\beta/2+\delta/2)} \cos \frac{\gamma}{2} \end{bmatrix} \\ &= e^{i\alpha} \begin{bmatrix} e^{-i\beta/2} & 0 \\ 0 & e^{i\beta/2} \end{bmatrix} \begin{bmatrix} \cos \gamma / 2 & -\sin \gamma / 2 \\ \sin \gamma / 2 & \cos \gamma / 2 \end{bmatrix} \begin{bmatrix} e^{-i\delta/2} & 0 \\ 0 & e^{i\delta/2} \end{bmatrix} \\ &= e^{i\alpha} R_z(\beta) R_y(\gamma) R_z(\delta) \end{aligned}$$

# Corollary

Set  $A$ ,  $B$ ,  $C$  as

$$A \equiv R_z(\beta)R_y\left(\frac{\gamma}{2}\right)$$

$$B \equiv R_y\left(-\frac{\gamma}{2}\right)R_z\left(-\frac{\delta + \beta}{2}\right)$$

$$C \equiv R_z\left(\frac{\delta - \beta}{2}\right)$$

Then

$$ABC = I, \quad U = e^{i\alpha} AXBXC$$

We will construct an arbitrary controlled- $U$  gate using  $A$ ,  $B$ , and  $C$

# Corollary

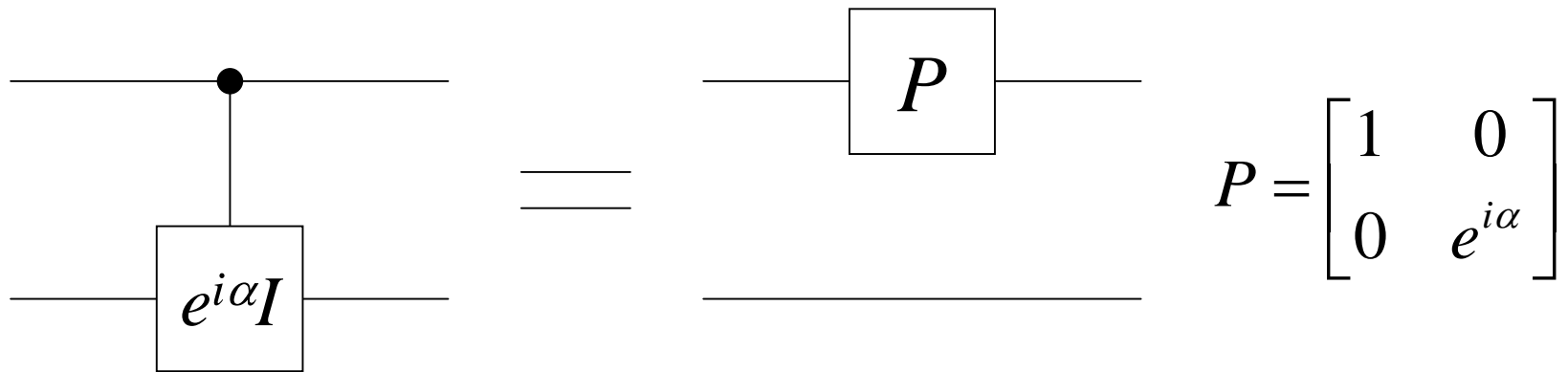
## Proof

$$ABC = R_z(\beta)R_y\left(\frac{\gamma}{2}\right)R_y\left(-\frac{\gamma}{2}\right)R_z\left(-\frac{\delta}{2}-\frac{\beta}{2}\right)R_z\left(\frac{\delta}{2}-\frac{\beta}{2}\right) = I$$

$$A = R_z(\beta)R_y\left(\frac{\gamma}{2}\right) \quad B = R_y\left(-\frac{\gamma}{2}\right)R_z\left(-\frac{\delta+\beta}{2}\right) \quad C = R_z\left(\frac{\delta-\beta}{2}\right)$$

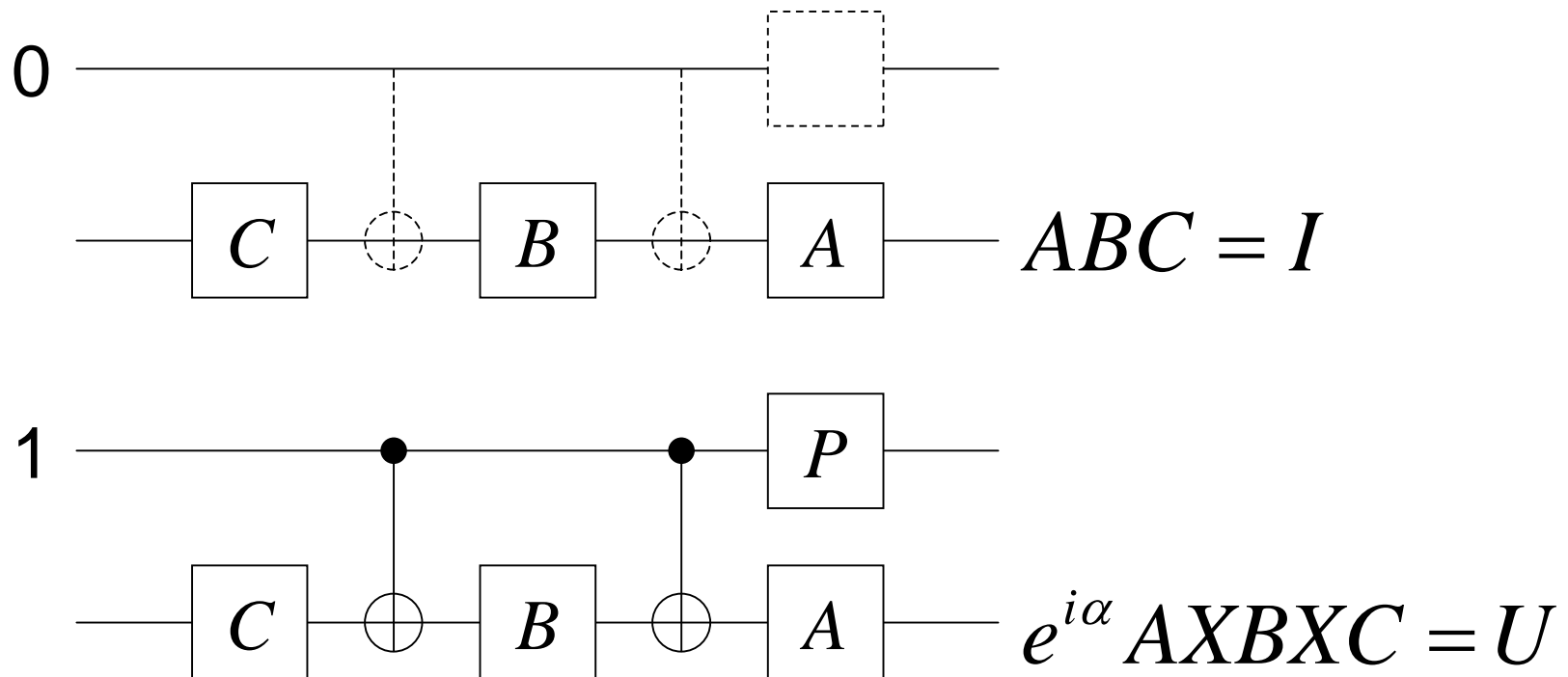
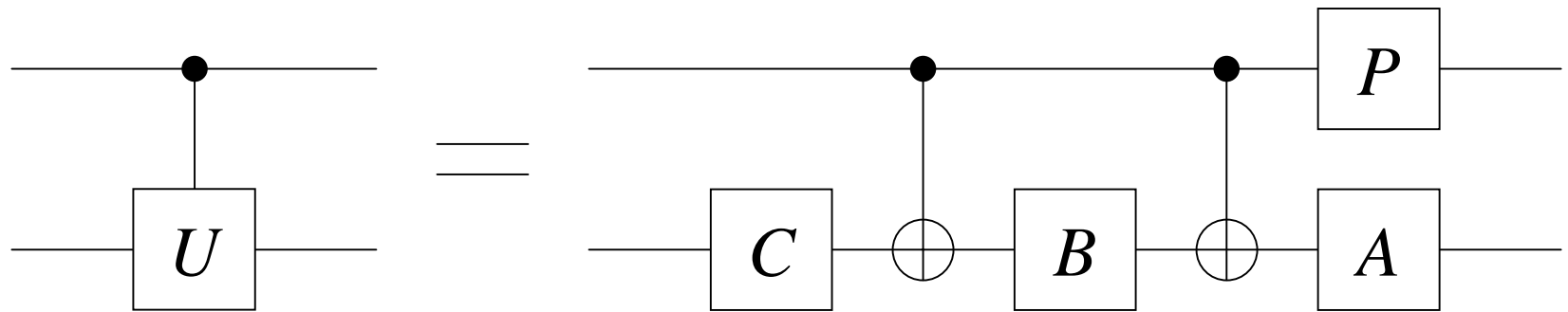
$$\begin{aligned} e^{i\alpha} AXBXC &= e^{i\alpha} AXR_y\left(-\frac{\gamma}{2}\right)XXR_z\left(-\frac{\delta+\beta}{2}\right)XC && \begin{array}{l} XX = I \\ XR_y(\theta)X = R_y(-\theta) \\ XR_z(\theta)X = R_z(-\theta) \end{array} \\ &= e^{i\alpha} AR_y\left(\frac{\gamma}{2}\right)R_z\left(\frac{\delta+\beta}{2}\right)C \\ &= e^{i\alpha} R_z(\beta)R_y\left(\frac{\gamma}{2}\right)R_y\left(\frac{\gamma}{2}\right)R_z\left(\frac{\delta+\beta}{2}\right)R_z\left(\frac{\delta-\beta}{2}\right) \\ &= e^{i\alpha} R_z(\beta)R_y(\gamma)R_z(\delta) = U \end{aligned}$$

# Phase shifter



$$\begin{array}{l} |00\rangle \rightarrow |00\rangle \\ |01\rangle \rightarrow |01\rangle \\ |10\rangle \rightarrow e^{i\alpha} |10\rangle \\ |11\rangle \rightarrow e^{i\alpha} |11\rangle \end{array} \iff P \otimes I = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & e^{i\alpha} & 0 \\ 0 & 0 & 0 & e^{i\alpha} \end{bmatrix}$$

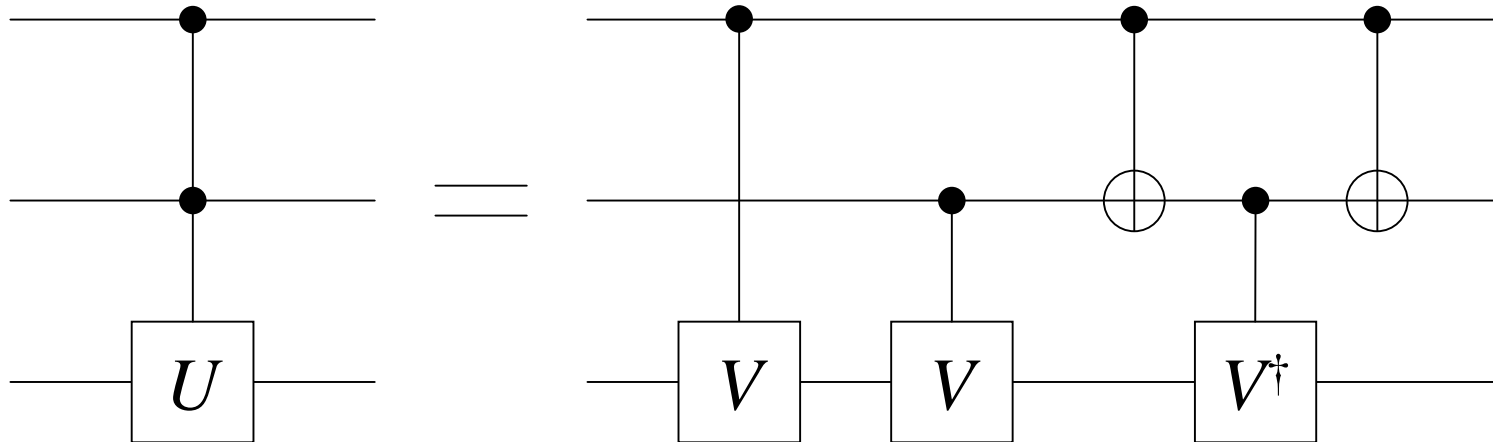
# Controlled- $U$ gate



# Road to universality proof

1. An arbitrary controlled- $U$  gate can be implemented using only single qubit gates and CNOT
2. An arbitrary (controlled) $^n$ - $U$  gate can be implemented using single qubit gates and CNOT
3. Two-level unitary gates are universal
4. Single-qubit gates and CNOT are universal
5. Hadamard,  $S$ ,  $T$ , and CNOT are universal

# (Controlled)<sup>2</sup>- $U$ gate



## Example

$$V^2 = U$$

$$VV^\dagger = I$$

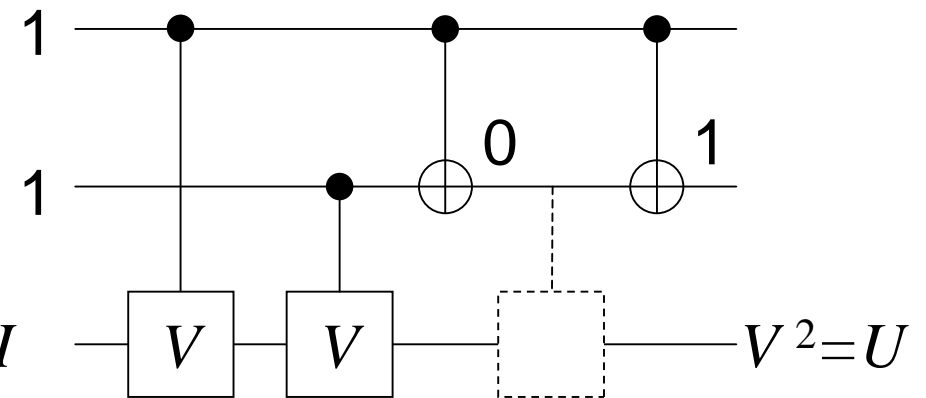
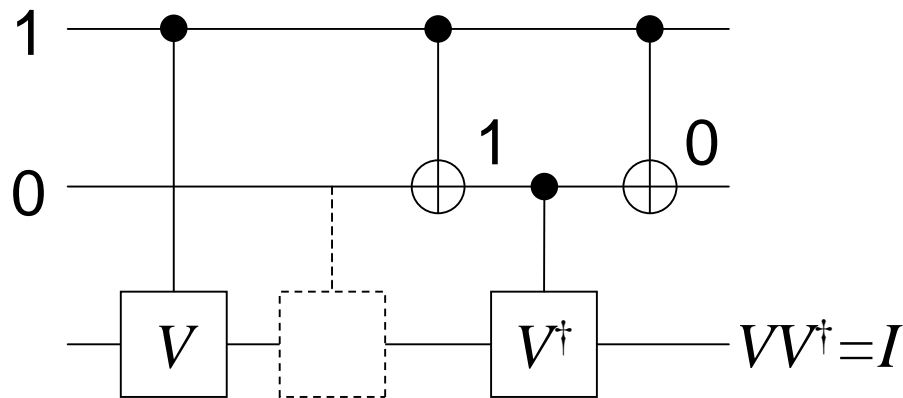
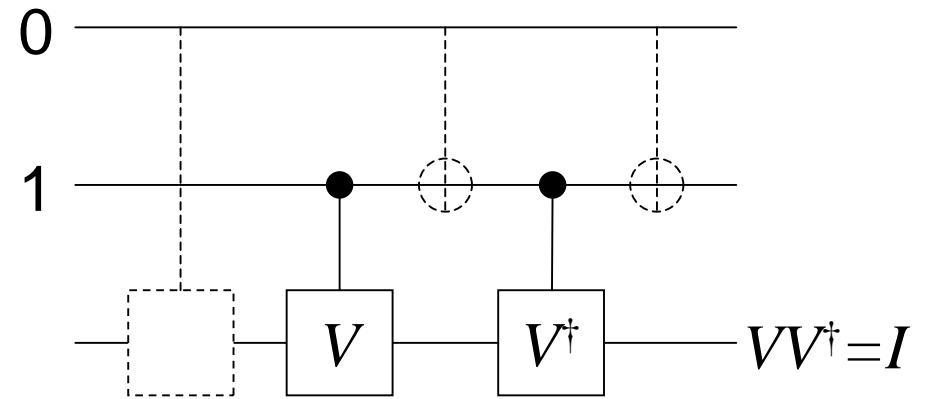
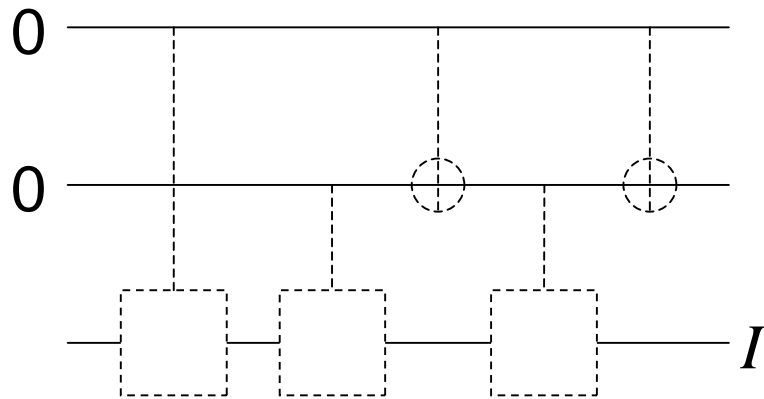
$$T^2 = S$$

$$(HSH)^2 = X$$

$$S^2 = Z \quad \left( \frac{I + iH}{\sqrt{2}} \right)^2 = iH$$



# (Controlled)<sup>2</sup>- $U$ gate



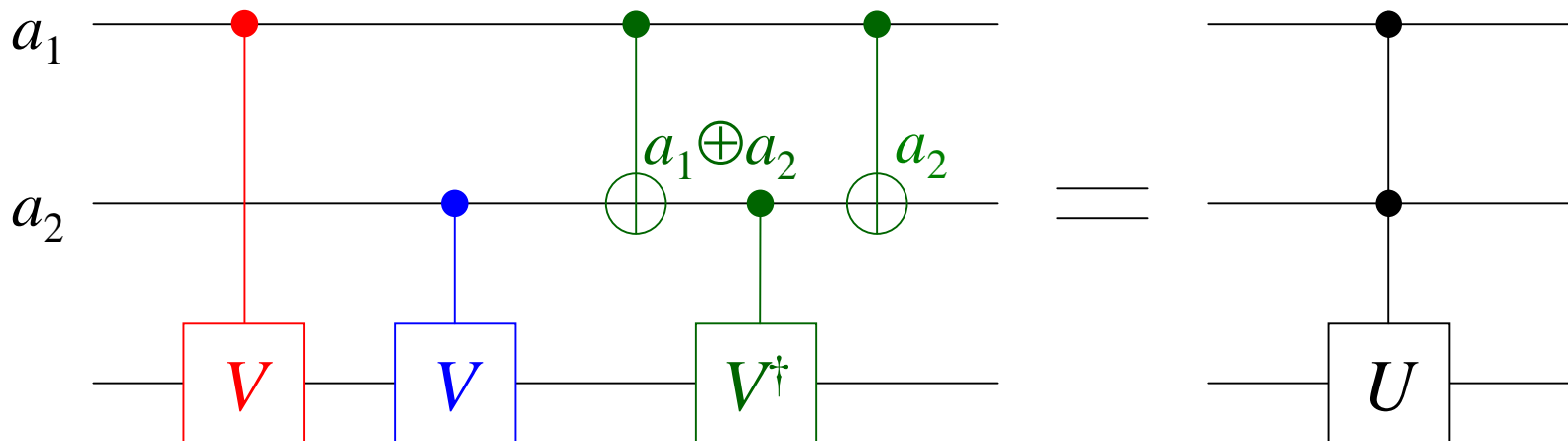
# (Controlled)<sup>2</sup>-*U* gate

$$a_1 + a_2 - (a_1 \oplus a_2) = 2 \times a_1 \cdot a_2$$

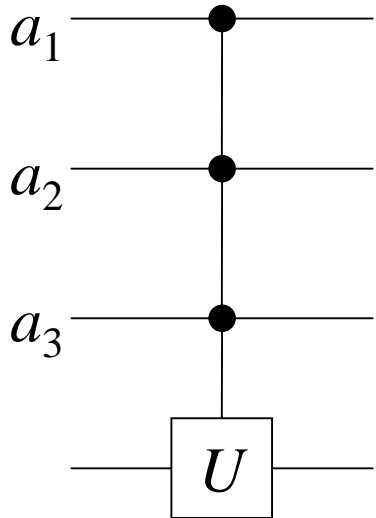
+, −, and × are the ordinary arithmetic operations

$a_1$	$a_2$	$a_1 \oplus a_2$	$2a_1 \cdot a_2$
0	0	0	0
0	1	1	0
1	0	1	0
1	1	0	2

$$V^{a_1} V^{a_2} (V^\dagger)^{a_1 \oplus a_2} = (V^2)^{a_1 \cdot a_2} = U^{a_1 \cdot a_2}$$

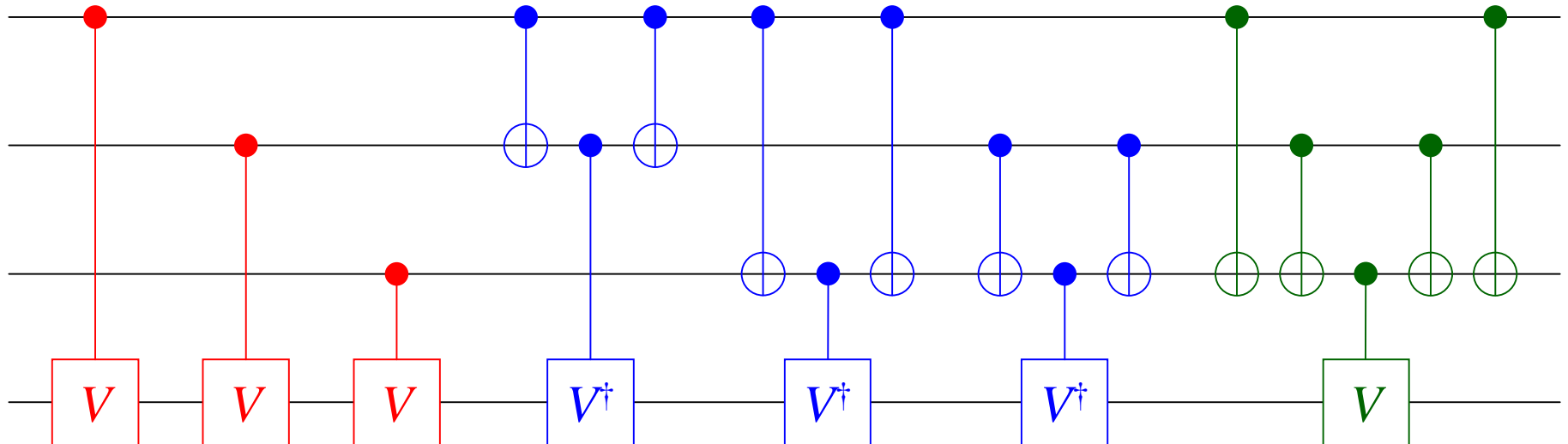


# (Controlled)<sup>3</sup>- $U$ gate



$$a_1 + a_2 + a_3 - (a_1 \oplus a_2) - (a_1 \oplus a_3) - (a_2 \oplus a_3) + (a_1 \oplus a_2 \oplus a_3) = 4 \times a_1 \cdot a_2 \cdot a_3$$

$$V^{a_1} V^{a_2} V^{a_3} (V^\dagger)^{a_1 \oplus a_2} (V^\dagger)^{a_1 \oplus a_3} (V^\dagger)^{a_2 \oplus a_3} V^{a_1 \oplus a_2 \oplus a_3} = (V^4)^{a_1 \cdot a_2 \cdot a_3} = U^{a_1 \cdot a_2 \cdot a_3}$$

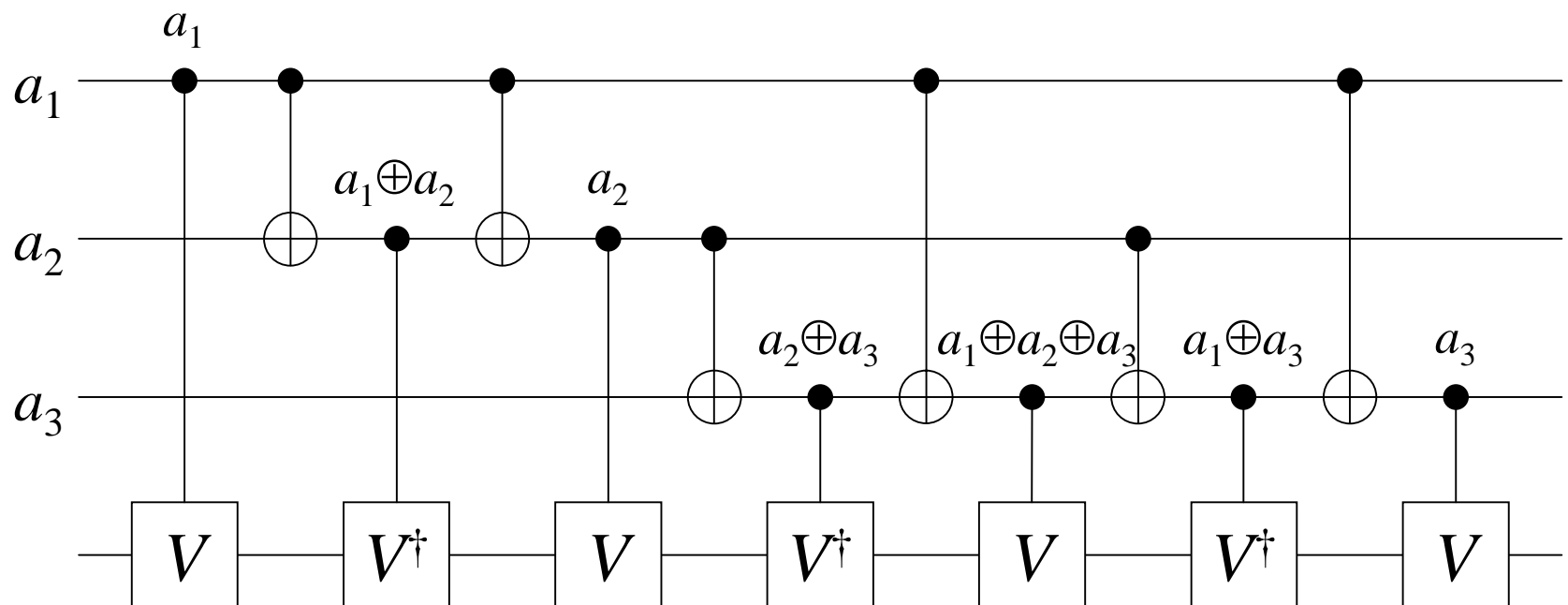


# Gate construction by a Gray code

$a_1$	1	1	0	0	1	1	0
$a_2$	0	1	1	1	1	0	0
$a_3$	0	0	0	1	1	1	1

Gray code; Only one bit changes from one entry to the next (patented by F. Gray)

$$V^{a_1} (V^\dagger)^{a_1 \oplus a_2} V^{a_2} (V^\dagger)^{a_2 \oplus a_3} V^{a_1 \oplus a_2 \oplus a_3} (V^\dagger)^{a_1 \oplus a_3} V^{a_3}$$



# (Controlled)<sup>n</sup>-U gate

Set  $V$  so that  $V^{2^{n-2}} = U$  and implement the identity

$$\sum_i a_i - \sum_{i < j} (a_i \oplus a_j) + \sum_{i < j < k} (a_i \oplus a_j \oplus a_k) - \dots + (-1)^{n-1} (a_1 \oplus a_2 \oplus \dots \oplus a_n)$$

$$= 2^{n-1} \times a_1 \cdot a_2 \cdot a_3 \cdots a_n$$

## Proof for $n = 3$

Can be proved for any  $n$  by induction

$$4 \times a_1 \cdot a_2 \cdot a_3$$

$$= 2(2a_1 \cdot a_2) \cdot a_3 \quad 2a_1 \cdot a_2 = a_1 + a_2 - (a_1 \oplus a_2)$$

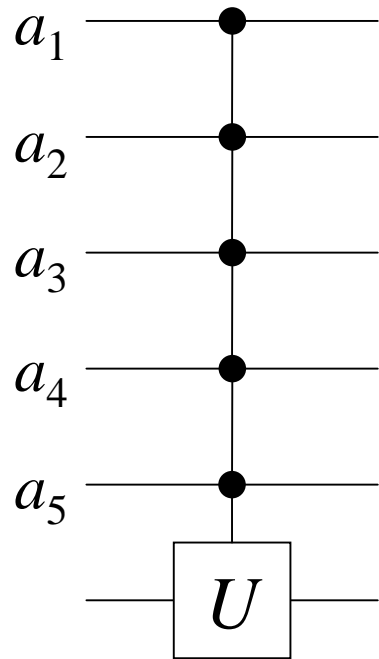
$$= 2[a_1 + a_2 - (a_1 \oplus a_2)] \cdot a_3$$

$$= 2a_1 \cdot a_3 + 2a_2 \cdot a_3 - 2(a_1 \oplus a_2) \cdot a_3$$

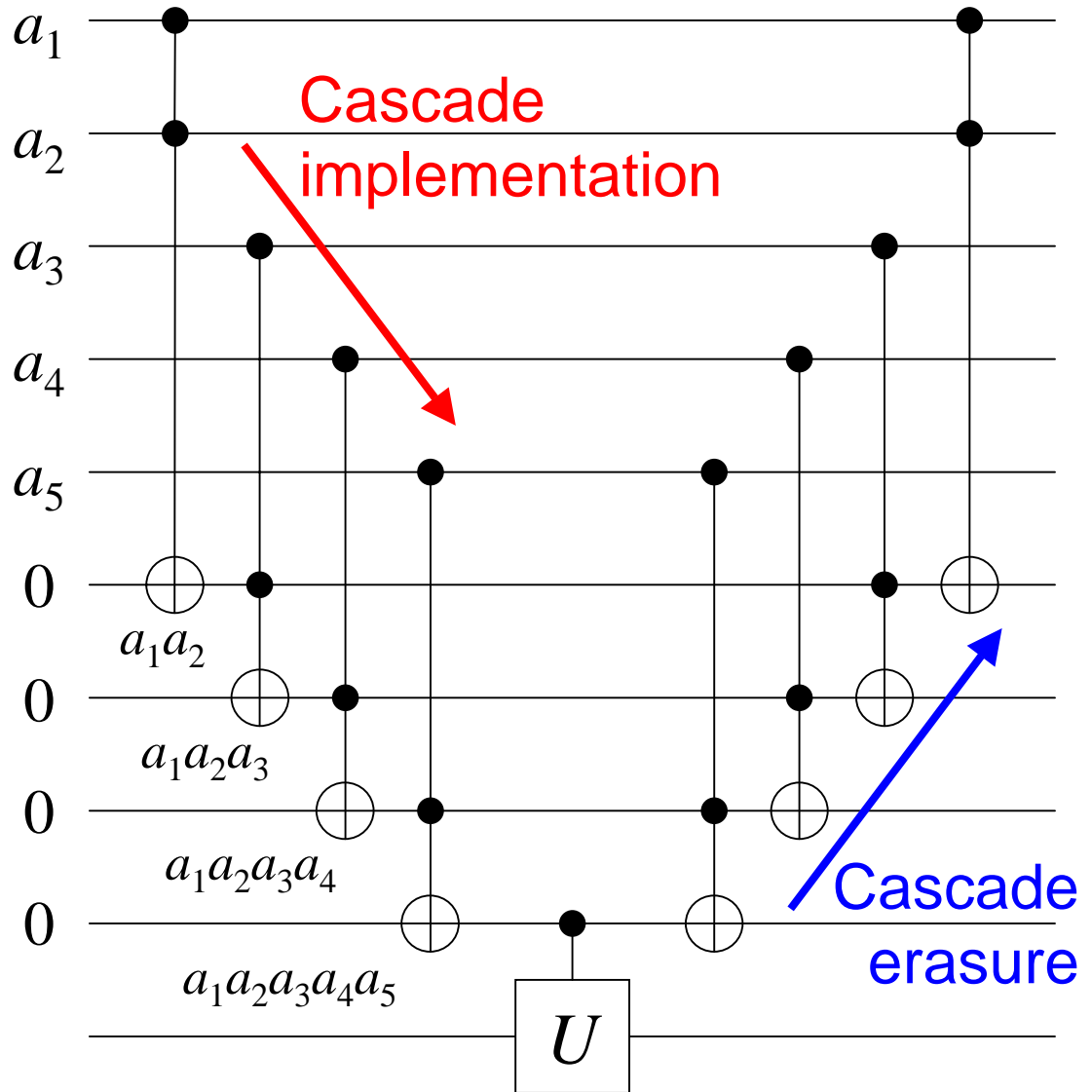
$$= [a_1 + a_3 - (a_1 \oplus a_3)] + [a_2 + a_3 - (a_2 \oplus a_3)] - [(a_1 \oplus a_2) + a_3 - (a_1 \oplus a_2 \oplus a_3)]$$

$$= a_1 + a_2 + a_3 - (a_1 \oplus a_2) - (a_1 \oplus a_3) - (a_2 \oplus a_3) + (a_1 \oplus a_2 \oplus a_3)$$

# (Controlled)<sup>n</sup>-*U* gate



=



- ✓  $n - 1$  ancilla
- ✓  $2(n - 1)$  Toffoli
- ✓ 1 Controlled-*U*

# Road to universality proof

1. An arbitrary controlled- $U$  gate can be implemented using only single qubit gates and CNOT
2. An arbitrary (controlled) $^n$ - $U$  gate can be implemented using single qubit gates and CNOT
3. **Two-level unitary gates are universal**
4. Single-qubit gates and CNOT are universal
5. Hadamard,  $S$ ,  $T$ , and CNOT are universal

# Two-level unitary gate

## Two-level unitary matrix

Unitary matrix which acts nontrivially only two-or-fewer vector components

$$\begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & \alpha & \gamma \\ 0 & 0 & \beta & \delta \end{bmatrix} \quad \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & \alpha & \gamma & 0 \\ 0 & \beta & \delta & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix}$$

## Universality; $3 \times 3$ case

Breaking  $U$  up into the product of two-level unitary matrices

$$I = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix}$$

$$U = \begin{bmatrix} a & d & g \\ b & e & h \\ c & f & j \end{bmatrix} \quad \rightarrow \quad U_3 U_2 U_1 U = I \quad \Leftrightarrow \quad U = U_1^\dagger U_2^\dagger U_3^\dagger$$



# Two-level unitary gates are universal

$$U = \begin{bmatrix} a & d & g \\ b & e & h \\ c & f & j \end{bmatrix}$$



$$U_1 U = \begin{bmatrix} a' & d' & g' \\ 0 & e' & h' \\ c' & f' & j' \end{bmatrix}$$




$$U_2 U_1 U = \begin{bmatrix} 1 & d'' & g'' \\ 0 & e'' & h'' \\ 0 & f'' & j'' \end{bmatrix}$$

$$U_1 \equiv \begin{cases} I & (b = 0) \\ \begin{bmatrix} \frac{a^*}{\sqrt{|a|^2+|b|^2}} & \frac{b^*}{\sqrt{|a|^2+|b|^2}} & 0 \\ b & -a & 0 \\ \frac{a}{\sqrt{|a|^2+|b|^2}} & \frac{b}{\sqrt{|a|^2+|b|^2}} & 1 \end{bmatrix} & (b \neq 0) \end{cases}$$

$$U_2 \equiv \begin{cases} \begin{bmatrix} a'^* & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix} & (c' = 0) \\ \begin{bmatrix} \frac{a'^*}{\sqrt{|a'|^2+|c'|^2}} & 0 & \frac{c'^*}{\sqrt{|a'|^2+|c'|^2}} \\ 0 & 1 & 0 \\ \frac{c'}{\sqrt{|a'|^2+|c'|^2}} & 0 & -a' \end{bmatrix} & (c' \neq 0) \end{cases}$$

# Two-level unitary gates are universal

$$U_2 U_1 U = \begin{bmatrix} 1 & d'' & g'' \\ 0 & e'' & h'' \\ 0 & f'' & j'' \end{bmatrix} = \begin{bmatrix} 1 & 0 & 0 \\ 0 & e'' & h'' \\ 0 & f'' & j'' \end{bmatrix}$$

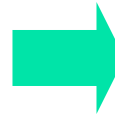


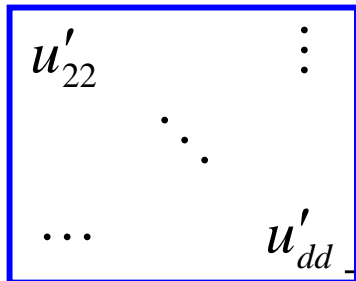
$$U_3 \equiv \begin{bmatrix} 1 & 0 & 0 \\ 0 & e''^* & f''^* \\ 0 & h''^* & j''^* \end{bmatrix}$$

$$U_3 U_2 U_1 U = I \iff U = U_1^\dagger U_2^\dagger U_3^\dagger$$

For  $d$ -dimensional  $U$ , we repeat this procedure

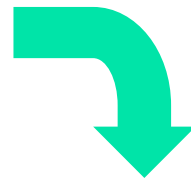
$$U = \begin{bmatrix} u_{11} & \cdots & u_{1d} \\ \vdots & \ddots & \vdots \\ u_{d1} & \cdots & u_{dd} \end{bmatrix} \xrightarrow{\quad} U_{d-1} U_{d-2} \cdots U_1 U = \begin{bmatrix} 1 & 0 & \cdots & 0 \\ 0 & u'_{22} & & \vdots \\ \vdots & & \ddots & \\ 0 & \cdots & & u'_{dd} \end{bmatrix}$$





# Two-level unitary gates are universal

$$U_{d-1}U_{d-2}\cdots U_1U = \begin{bmatrix} 1 & 0 & \cdots & 0 \\ 0 & u'_{22} & & \vdots \\ \vdots & & \ddots & \\ 0 & \cdots & & u'_{dd} \end{bmatrix}$$



$$\underbrace{U_{2d-3}\cdots U_d}_{d-2} \underbrace{(U_{d-1}U_{d-2}\cdots U_1U)}_{d-1} = \begin{bmatrix} 1 & 0 & 0 & \cdots & 0 \\ 0 & 1 & 0 & \cdots & 0 \\ 0 & 0 & u''_{33} & \cdots & \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & & \cdots & u''_{dd} \end{bmatrix}$$



$$\underbrace{U_k \cdots U_1 U = I}_{k \leq (d-1) + (d-2) + \cdots + 1 = \frac{d(d-1)}{2}} \Leftrightarrow U = U_1^\dagger \cdots U_k^\dagger$$

# Road to universality proof

1. An arbitrary controlled- $U$  gate can be implemented using only single qubit gates and CNOT
2. An arbitrary (controlled) $^n$ - $U$  gate can be implemented using single qubit gates and CNOT
3. Two-level unitary gates are universal
4. **Single-qubit gates and CNOT are universal**
5. Hadamard,  $S$ ,  $T$ , and CNOT are universal

# Single qubit gates & CNOT are universal

## Strategy

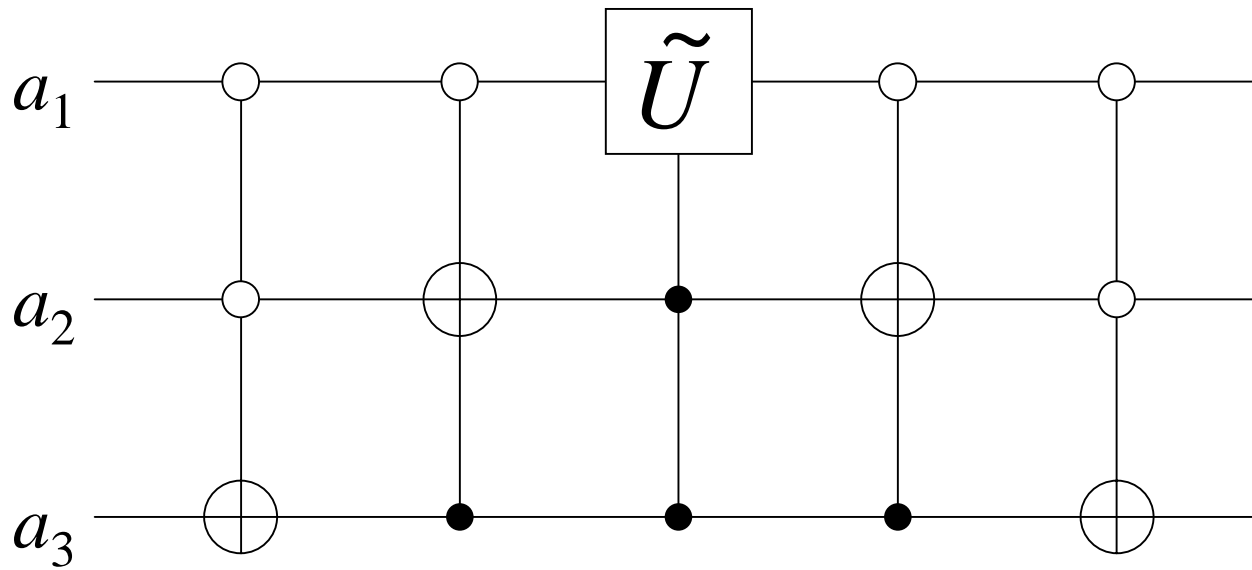
To show that single qubit & CNOT gates can implement an arbitrary two-level unitary matrix

## Example

$8 \times 8$   $U$  acting nontrivially only on  $|000\rangle$  and  $|111\rangle$

$$U = \begin{bmatrix} \alpha & 0 & 0 & 0 & 0 & 0 & 0 & \gamma \\ 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 \\ \beta & 0 & 0 & 0 & 0 & 0 & 0 & \delta \end{bmatrix} \quad \longrightarrow \quad U \begin{bmatrix} a \\ b \\ c \\ d \\ e \\ f \\ g \\ h \end{bmatrix} = \begin{bmatrix} \alpha a + \gamma h \\ b \\ c \\ d \\ e \\ f \\ g \\ \beta a + \delta h \end{bmatrix}$$

# Single qubit gates & CNOT are universal

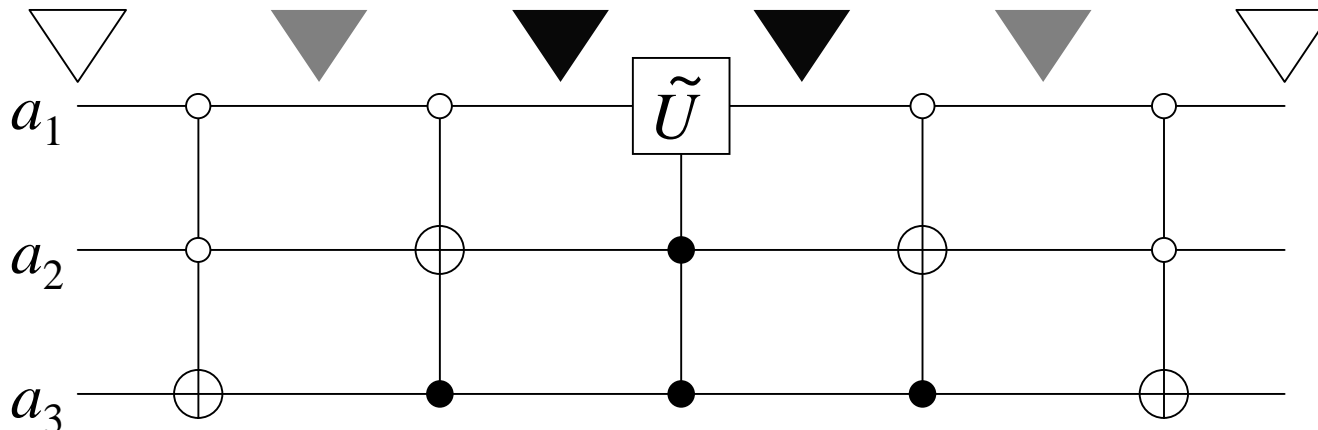


$$\tilde{U} = \begin{bmatrix} \alpha & \gamma \\ \beta & \delta \end{bmatrix}$$

$a_1$	0	0	0	1
$a_2$	0	0	1	1
$a_3$	0	1	1	1

We want to apply a controlled gate with the target bit  $|a_1\rangle$

# Single qubit gates & CNOT are universal



$$\tilde{U} = \begin{bmatrix} \alpha & \gamma \\ \beta & \delta \end{bmatrix}$$

$a_1 a_2 a_3$	Stage 1	Stage 2	Stage 3	Stage 4	Stage 5
000	$a$	$b$	$b$	$b$	$\alpha a + \gamma h$
001	$b$	$a$	$d$	$d$	$b$
010	$c$	$c$	$c$	$c$	$c$
011	$d$	$d$	$a$	$\alpha a + \gamma h$	$d$
100	$e$	$e$	$e$	$e$	$e$
101	$f$	$f$	$f$	$f$	$f$
110	$g$	$g$	$g$	$g$	$g$
111	$h$	$h$	$h$	$\beta a + \delta h$	$\beta a + \delta h$

# Road to universality proof

1. An arbitrary controlled- $U$  gate can be implemented using only single qubit gates and CNOT
2. An arbitrary (controlled) $^n$ - $U$  gate can be implemented using single qubit gates and CNOT
3. Two-level unitary gates are universal
4. Single-qubit gates and CNOT are universal
5. Hadamard,  $S$ ,  $T$ , and CNOT are universal



# Discrete set of universal gates

## Why a discrete set of gates?

It can be used to perform quantum computation in an ***error-resistant*** fashion

## Problem

The set of unitary operations is ***continuous***

## Strategy

To show that a discrete set can be used to ***approximate*** any unitary operation to an ***arbitrary accuracy***

# Approximation by $H$ & $T$

$$T = \begin{bmatrix} 1 & 0 \\ 0 & e^{i\pi/4} \end{bmatrix} = e^{i\pi/8} \begin{bmatrix} e^{-i\pi/8} & 0 \\ 0 & e^{i\pi/8} \end{bmatrix} \cong R_z(\pi/4)$$

$$HTH = e^{i\pi/8} HR_z(\pi/4)H \cong R_x(\pi/4)$$



$$HZH = X \Rightarrow HR_z(\theta)H = R_x(\theta)$$

$$R_z(\pi/4)R_x(\pi/4) = R_{\hat{n}}(\theta)$$

$$\cos(\theta/2) \equiv \cos^2(\pi/8)$$

$\theta$ ; irrational multiple of  $2\pi$

$$\sin(\theta/2) = \sqrt{1 - \cos^4(\pi/8)} = \sin(\pi/8)\sqrt{1 + \cos^2(\pi/8)}$$

$$\hat{n} = \begin{bmatrix} n_x \\ n_y \\ n_z \end{bmatrix} \equiv \frac{1}{\sqrt{1 + \cos^2(\pi/8)}} \begin{bmatrix} \cos(\pi/8) \\ \sin(\pi/8) \\ \cos(\pi/8) \end{bmatrix}$$

# Approximation by $H$ & $T$

$$\begin{aligned} R_z\left(\frac{\pi}{4}\right)R_x\left(\frac{\pi}{4}\right) &= \exp\left(-i\frac{\pi}{8}Z\right)\exp\left(-i\frac{\pi}{8}X\right) && -iZX = Y \\ &= \left[\cos\frac{\pi}{8}I - i\sin\frac{\pi}{8}Z\right]\left[\cos\frac{\pi}{8}I - i\sin\frac{\pi}{8}X\right] \\ &= \cos^2\frac{\pi}{8}I - i\sin\frac{\pi}{8}\left[\cos\frac{\pi}{8}X + \sin\frac{\pi}{8}Y + \cos\frac{\pi}{8}Z\right] \\ &= \cos\frac{\theta}{2}I - i\sin\frac{\theta}{2}\left[n_xX + n_yY + n_zZ\right] = R_{\hat{n}}(\theta) \end{aligned}$$

$$\cos(\theta/2) \equiv \cos^2(\pi/8)$$

$$\sin(\theta/2) = \sqrt{1 - \cos^4(\pi/8)} = \sin(\pi/8)\sqrt{1 + \cos^2(\pi/8)}$$

$$\hat{n} = \begin{bmatrix} n_x \\ n_y \\ n_z \end{bmatrix} \equiv \frac{1}{\sqrt{1 + \cos^2(\pi/8)}} \begin{bmatrix} \cos(\pi/8) \\ \sin(\pi/8) \\ \cos(\pi/8) \end{bmatrix}$$

# Approximation by $H$ & $T$

## Weyl's theorem on uniform distribution

Let  $p$  be irrational, then the sequence  $\{p, 2p, 3p, \dots\}$  is uniformly distributed modulo 1

$n = 1$



$$\{n \theta / 2\pi \pmod{1}\}$$

$n = 10$



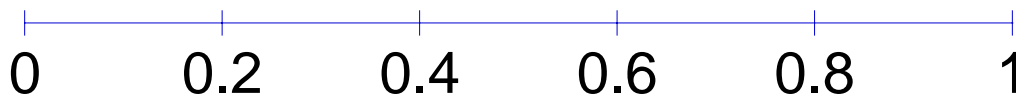
$n = 50$



$n = 100$



$n = 500$



The approximation to accuracy  $\varepsilon$  is realized through  $O(1/\varepsilon)$  times iterations

$$R_{\hat{n}}^{O(1/\varepsilon)}(\theta) \approx R_{\hat{n}}(\alpha)$$

# $H + S + T + \text{CNOT}$ are universal

$$HR_{\hat{n}}(\alpha)H = R_{\hat{m}}(\alpha)$$



$$U = R_{\hat{n}}(\beta)R_{\hat{m}}(\gamma)R_{\hat{n}}(\delta)$$



$$U \approx R_{\hat{n}}^{n_1}(\theta)R_{\hat{m}}^{n_2}(\theta)R_{\hat{n}}^{n_3}(\theta)$$

$$\begin{aligned} & n_x HXH + n_y HYH + n_z HZH \\ &= n_x X - n_y Y + n_z Z \end{aligned}$$

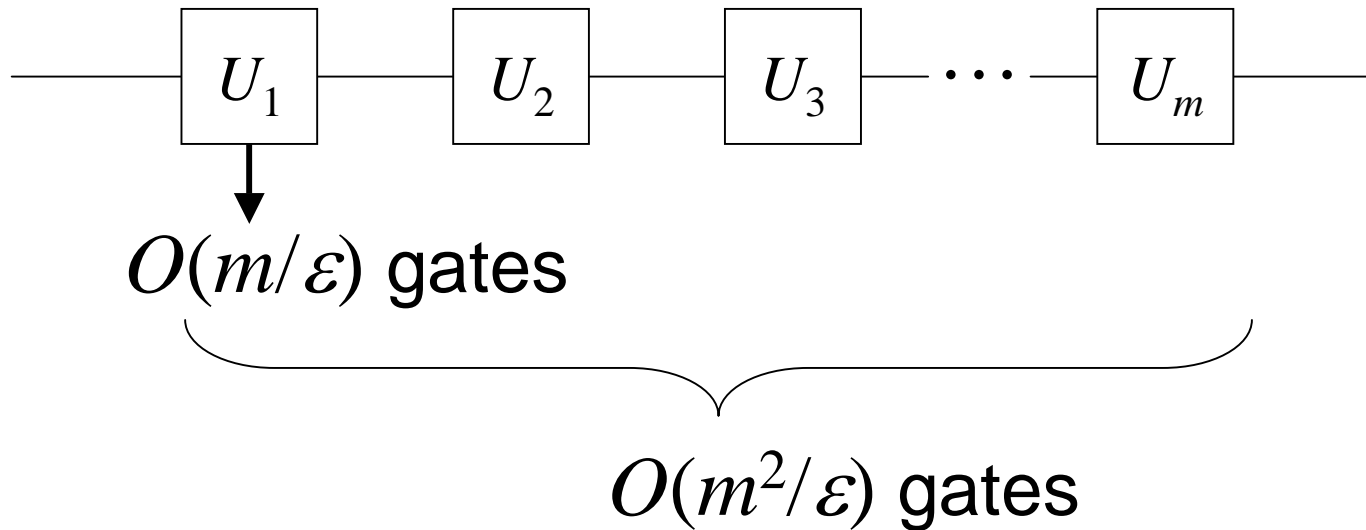
$$\hat{m} \equiv \frac{1}{\sqrt{1 + \cos^2(\pi/8)}} \begin{bmatrix} \cos(\pi/8) \\ -\sin(\pi/8) \\ \cos(\pi/8) \end{bmatrix}$$

$S$  has its own role in doing the approximation in a fault-tolerant fashion

***Is this construction efficient?***

# Efficiency

*Total accuracy  $\varepsilon$*

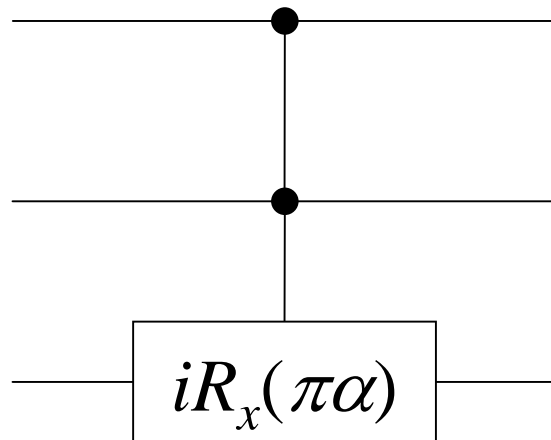


Solovay-Kitaev theorem

➡  $O(m \log^c (m / \varepsilon)) \quad c \approx 2$

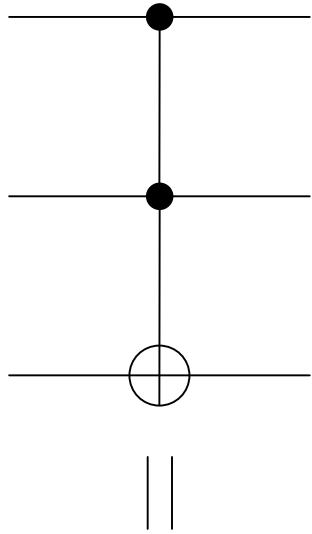
# Discrete set of universal gates

- $H$ ,  $S$ ,  $T$ , and CNOT
- $H$ ,  $S$ , CNOT, and Toffoli
- Deutsch gate

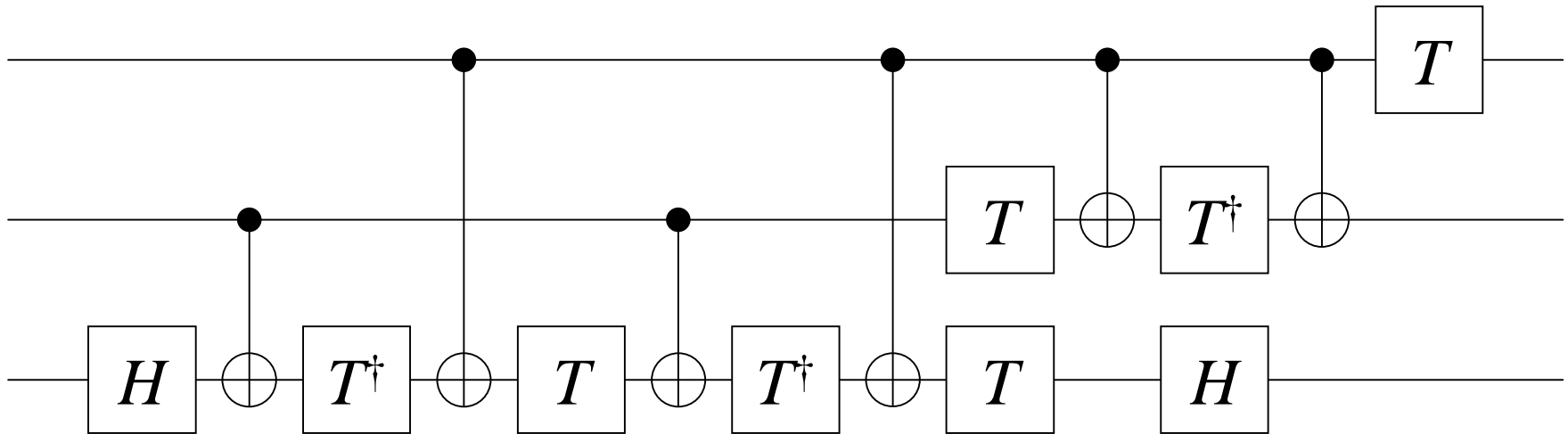


$\alpha$  ; irrational

# Quiz

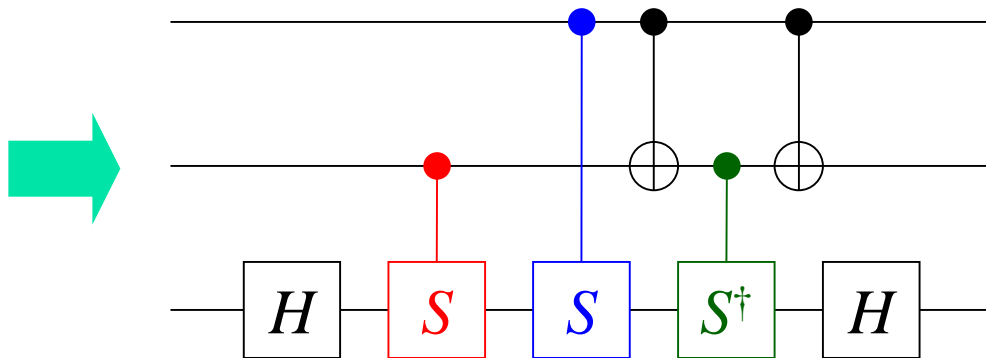
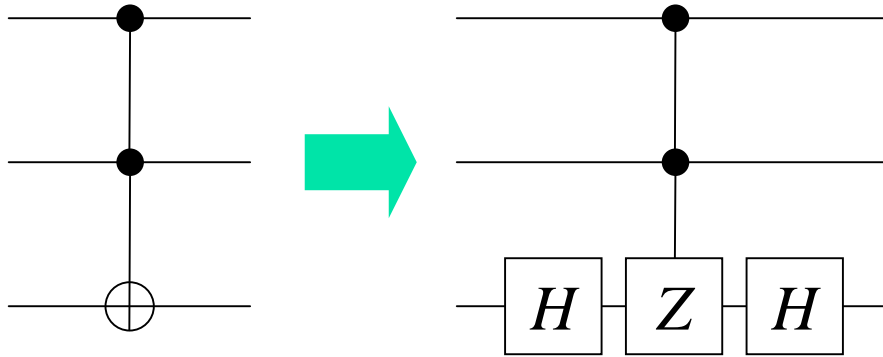


Construct Toffoli using only  $H$ ,  $T$ , and CNOT





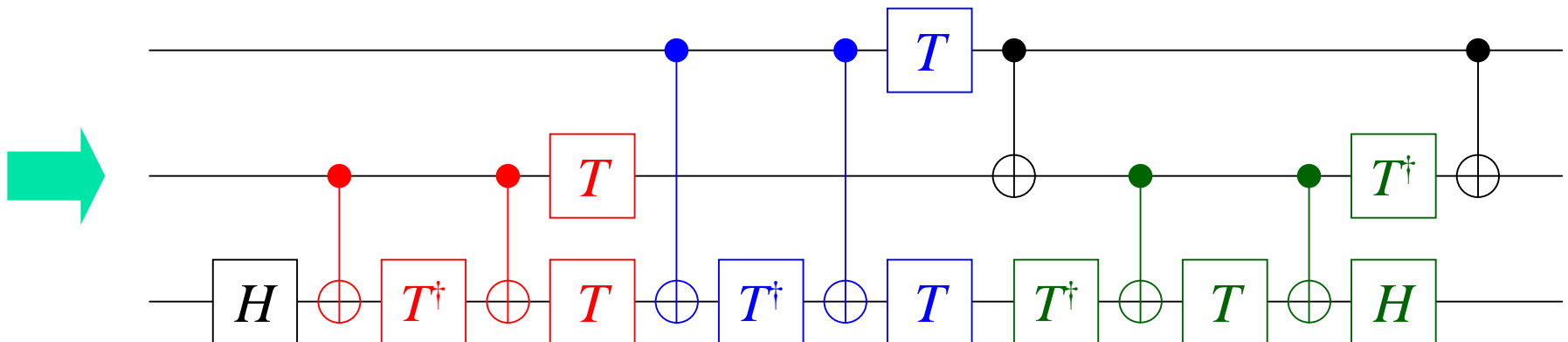
# Answer



$$S^2 = Z$$

$$S = e^{i\pi/4} T X T^\dagger X$$

$$S^\dagger = e^{-i\pi/4} X T X T^\dagger$$





# Answer

