

Grover's algorithm

School on Quantum Computing @Yagami

Day 1, Lesson 4

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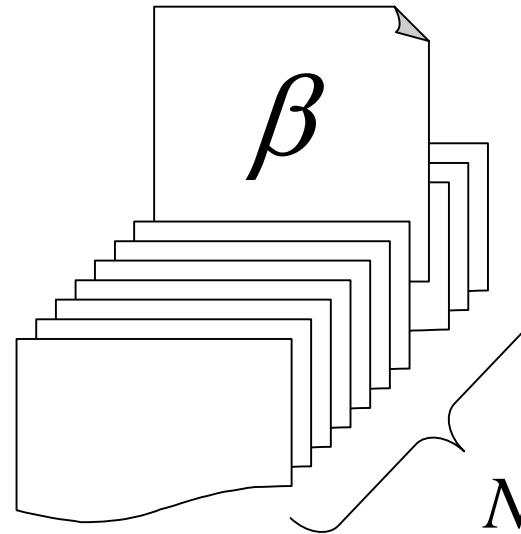
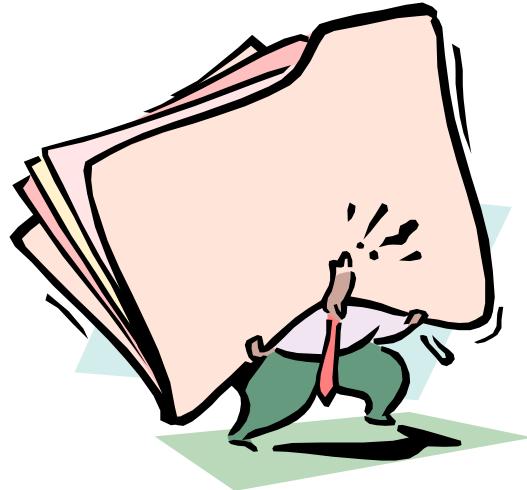


Outline

- Grover's algorithm
 - Oracle
 - Grover operator
 - Examples
 - 2-bit
 - 3bit
 - Geometric visualization
 - Performance and optimality

Database searching

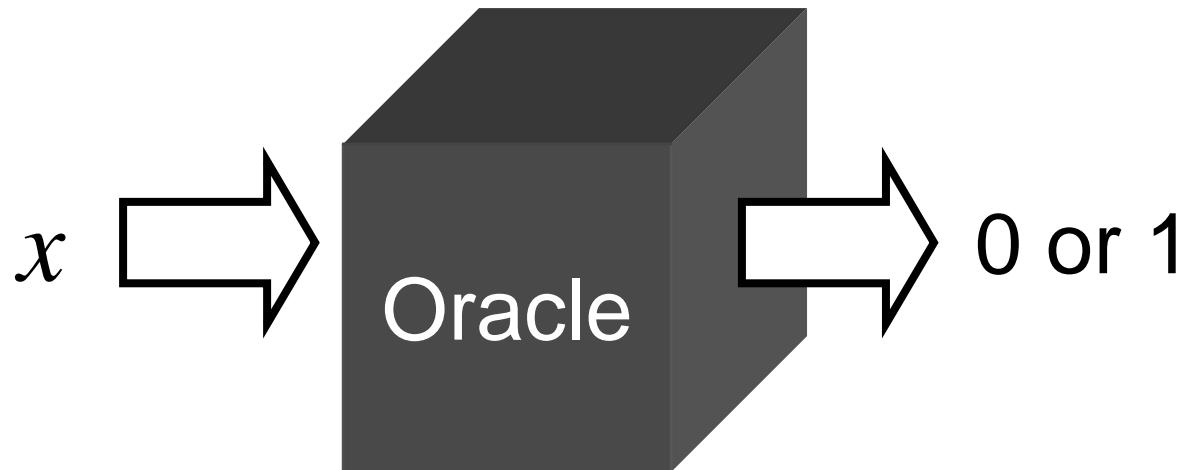
Find the desired file indexed as “ β ”
among $N = 2^n$ files



Oracle

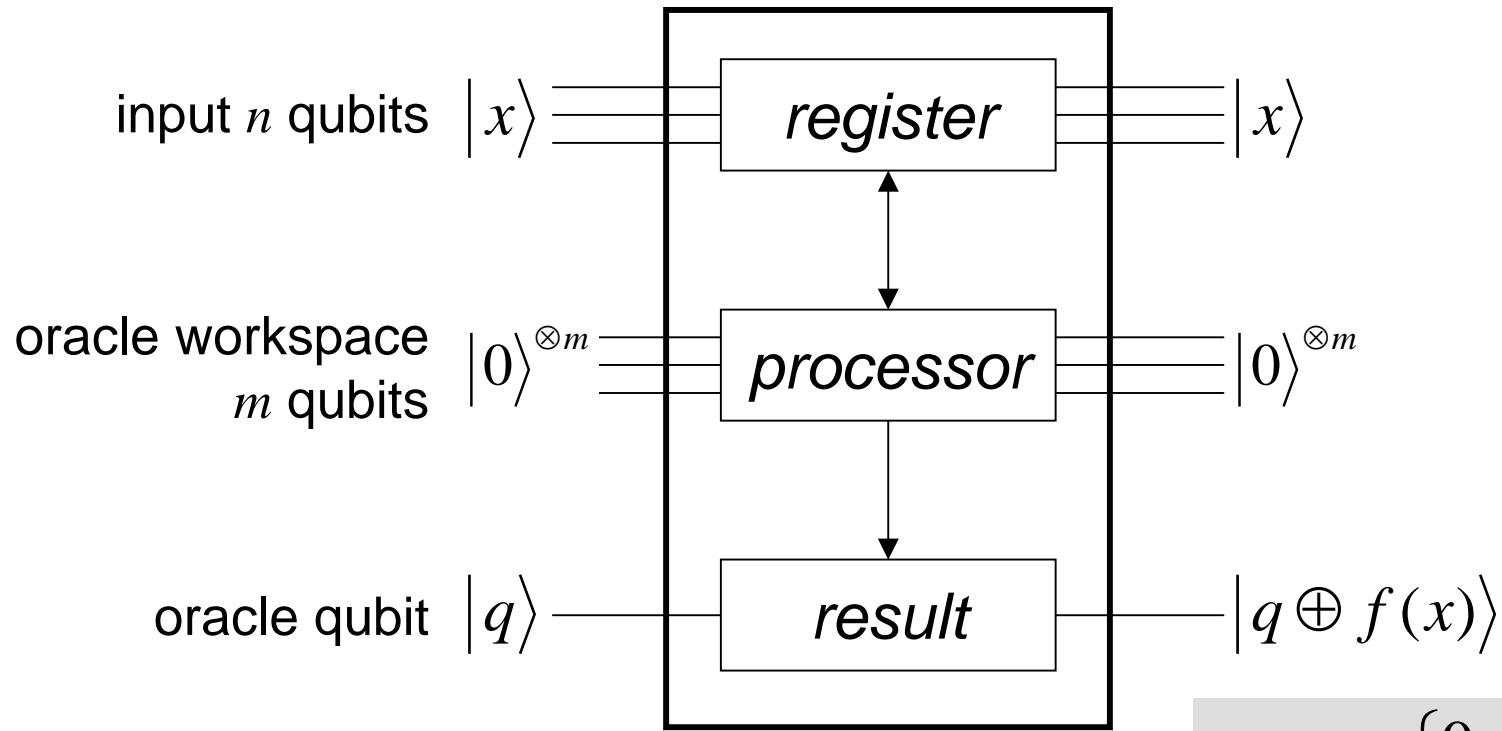
A black box that can *recognize* the solution, whose internal working is represented by a binary function $f(x)$

$$f(x) = \begin{cases} 0 & (x \neq \beta) \\ 1 & (x = \beta) \end{cases}$$



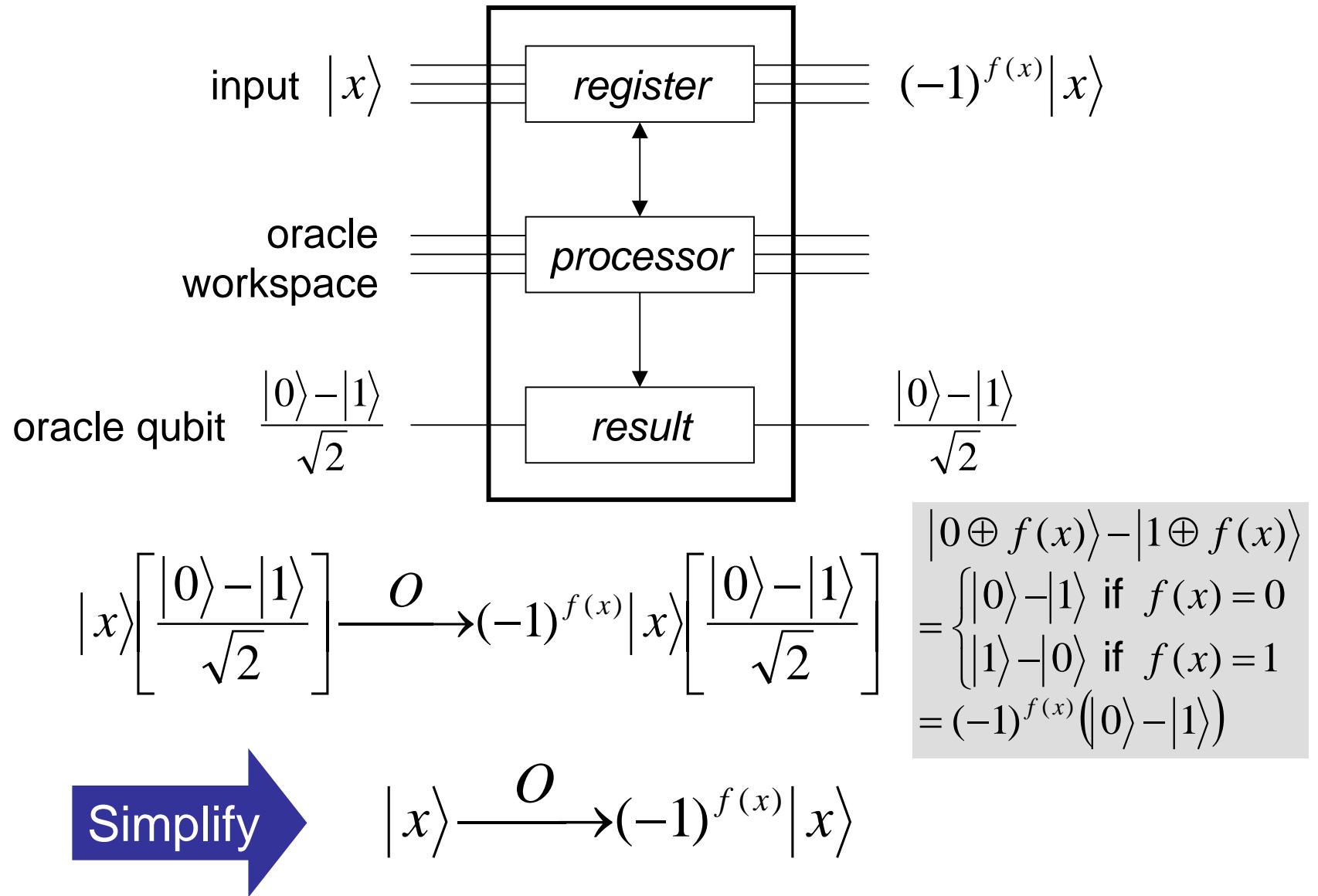
Oracle

$$|x\rangle|q\rangle \xrightarrow{O} |x\rangle|q \oplus f(x)\rangle$$

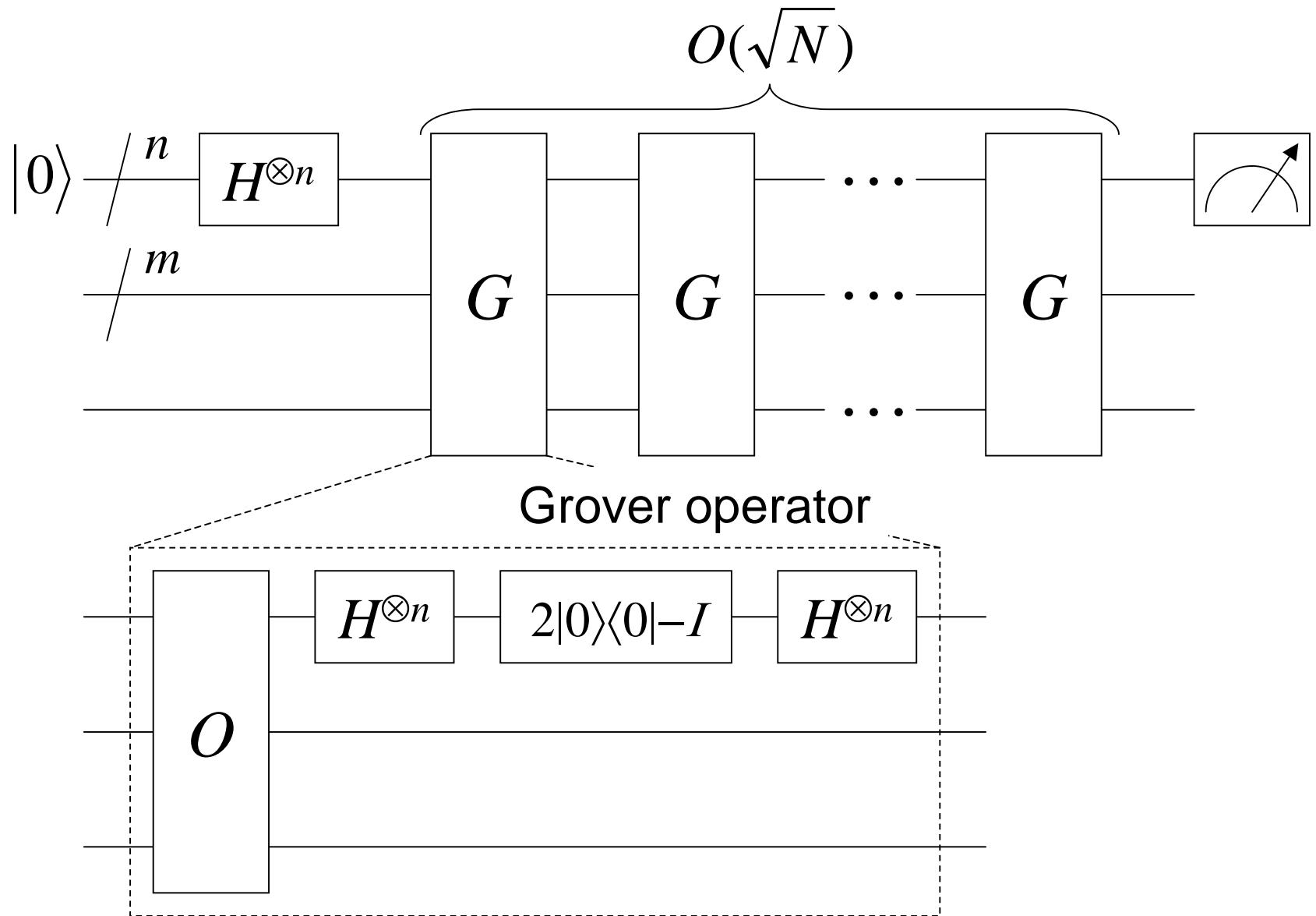


$$f(x) = \begin{cases} 0 & (x \neq \beta) \\ 1 & (x = \beta) \end{cases}$$

Oracle



Quantum search algorithm



Grover operator

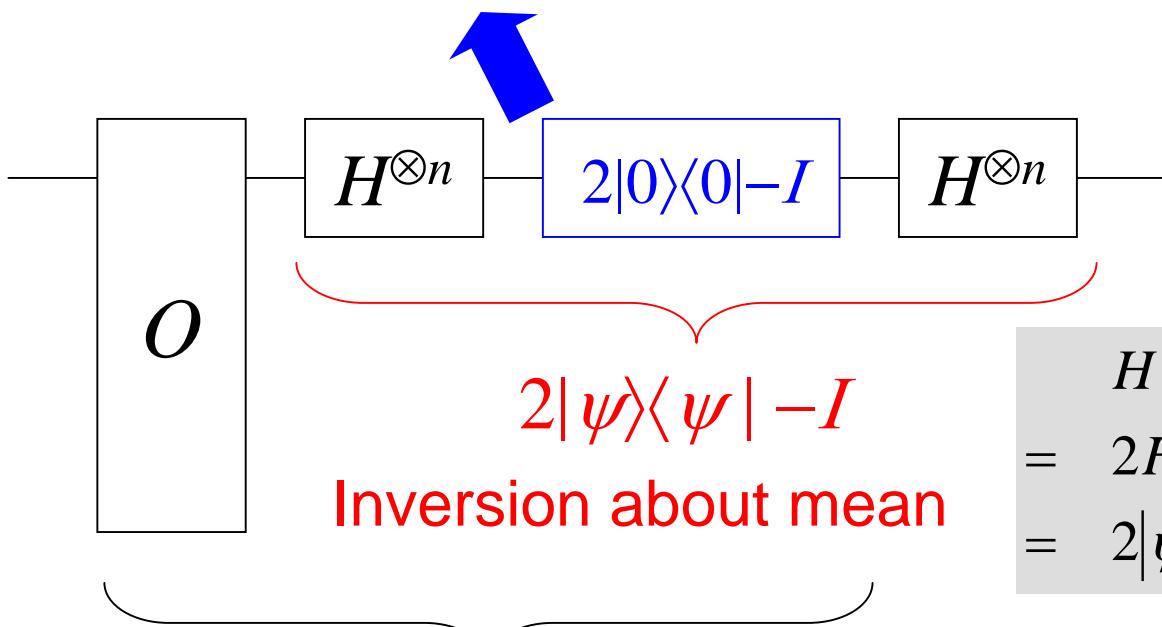
Phase shift operator

$$|0\rangle \rightarrow |0\rangle$$

$$|x\rangle \rightarrow -|x\rangle \quad (x \neq 0)$$

$$(2|0\rangle\langle 0| - I)|0\rangle = 2|0\rangle\langle 0|0\rangle - |0\rangle = |0\rangle$$

$$(2|0\rangle\langle 0| - I)|x\rangle = 2|0\rangle\langle 0|x\rangle - |x\rangle = -|x\rangle$$



$$G = (2|\psi\rangle\langle\psi| - I)O$$

$$\begin{aligned} & H^{\otimes n}(2|0\rangle\langle 0| - I)H^{\otimes n} \\ = & 2H^{\otimes n}|0\rangle\langle 0|H^{\otimes n} - H^{\otimes n}H^{\otimes n} \\ = & 2|\psi\rangle\langle\psi| - I \end{aligned}$$

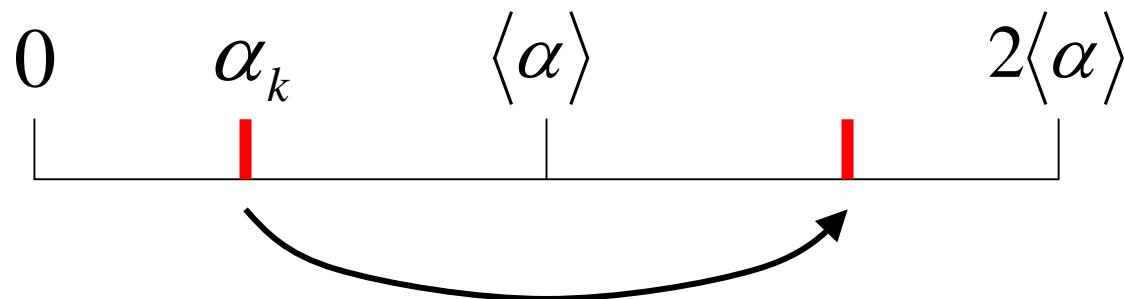
$$|\psi\rangle = \frac{1}{N^{1/2}} \sum_{x=0}^{N-1} |x\rangle$$

Grover operator

$$\begin{aligned}& (2|\psi\rangle\langle\psi| - I) \sum_k \underline{\alpha_k} |k\rangle \\&= 2N^{-1} \sum_{k,k',k''} \alpha_k |k'\rangle\langle k''|k\rangle - \sum_k \alpha_k |k\rangle \\&= 2N^{-1} \sum_{k,k'} \alpha_k |k'\rangle - \sum_k \alpha_k |k\rangle \\&= \sum_k \underline{(2\langle\alpha\rangle - \alpha_k)} |k\rangle\end{aligned}$$

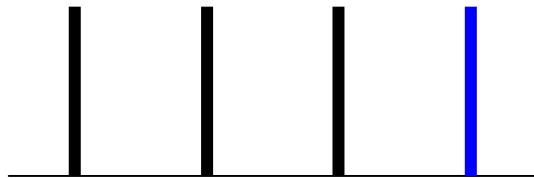
$$\begin{aligned}|\psi\rangle &= N^{-1/2} \sum |k'\rangle \\ \langle\psi| &= N^{-1/2} \sum \langle k''| \\ \langle k''|k\rangle &= \delta_{kk''} \\ \langle\alpha\rangle &= N^{-1} \sum \alpha_k \\ k' &\rightarrow k\end{aligned}$$

Inversion about mean

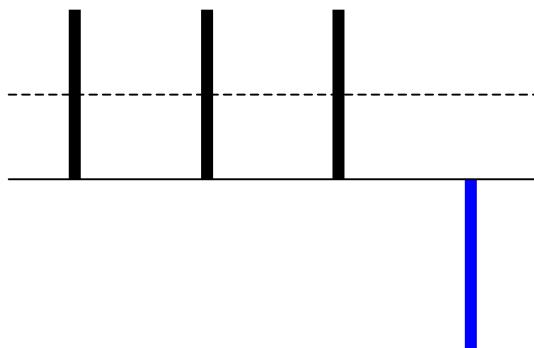


2-bit Grover

$|00\rangle \ |01\rangle \ |10\rangle \ |11\rangle$



$$|\psi\rangle = \frac{1}{2}|00\rangle + \frac{1}{2}|01\rangle + \frac{1}{2}|10\rangle + \frac{1}{2}|11\rangle$$



$$O|\psi\rangle = \frac{1}{2}|00\rangle + \frac{1}{2}|01\rangle + \frac{1}{2}|10\rangle - \frac{1}{2}|11\rangle$$



$$G|\psi\rangle = (2|\psi\rangle\langle\psi| - I)O|\psi\rangle = |11\rangle$$

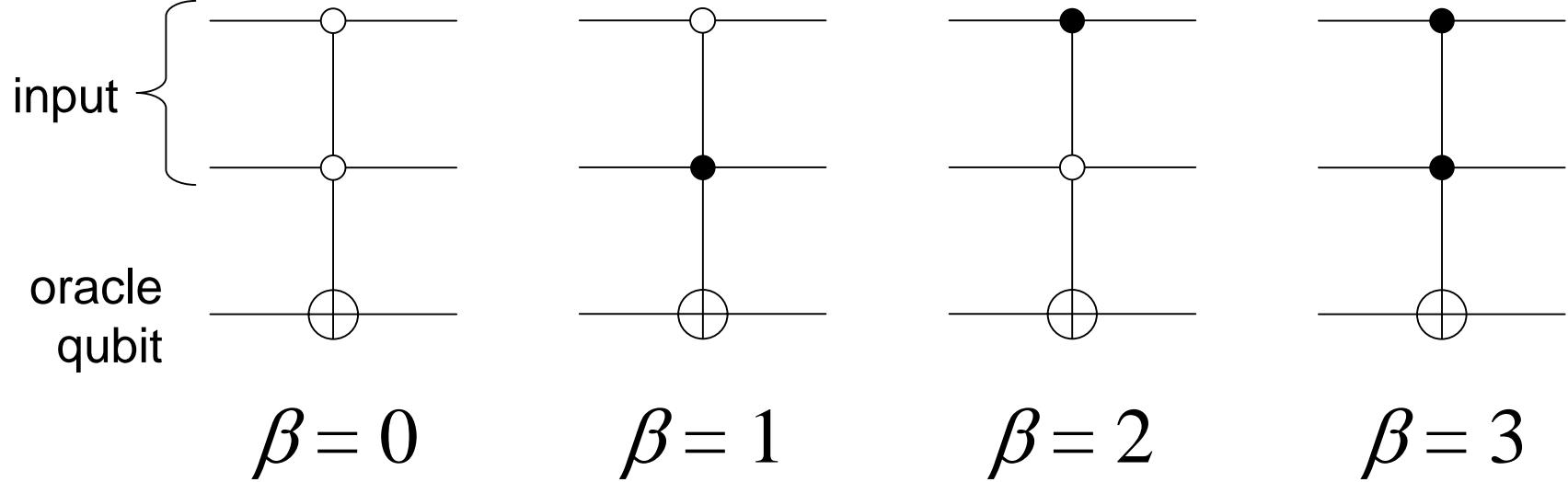
Observe now!



$$\langle\alpha\rangle = \frac{1}{4}\left(3 \times \frac{1}{2} - \frac{1}{2}\right) = \frac{1}{4}$$

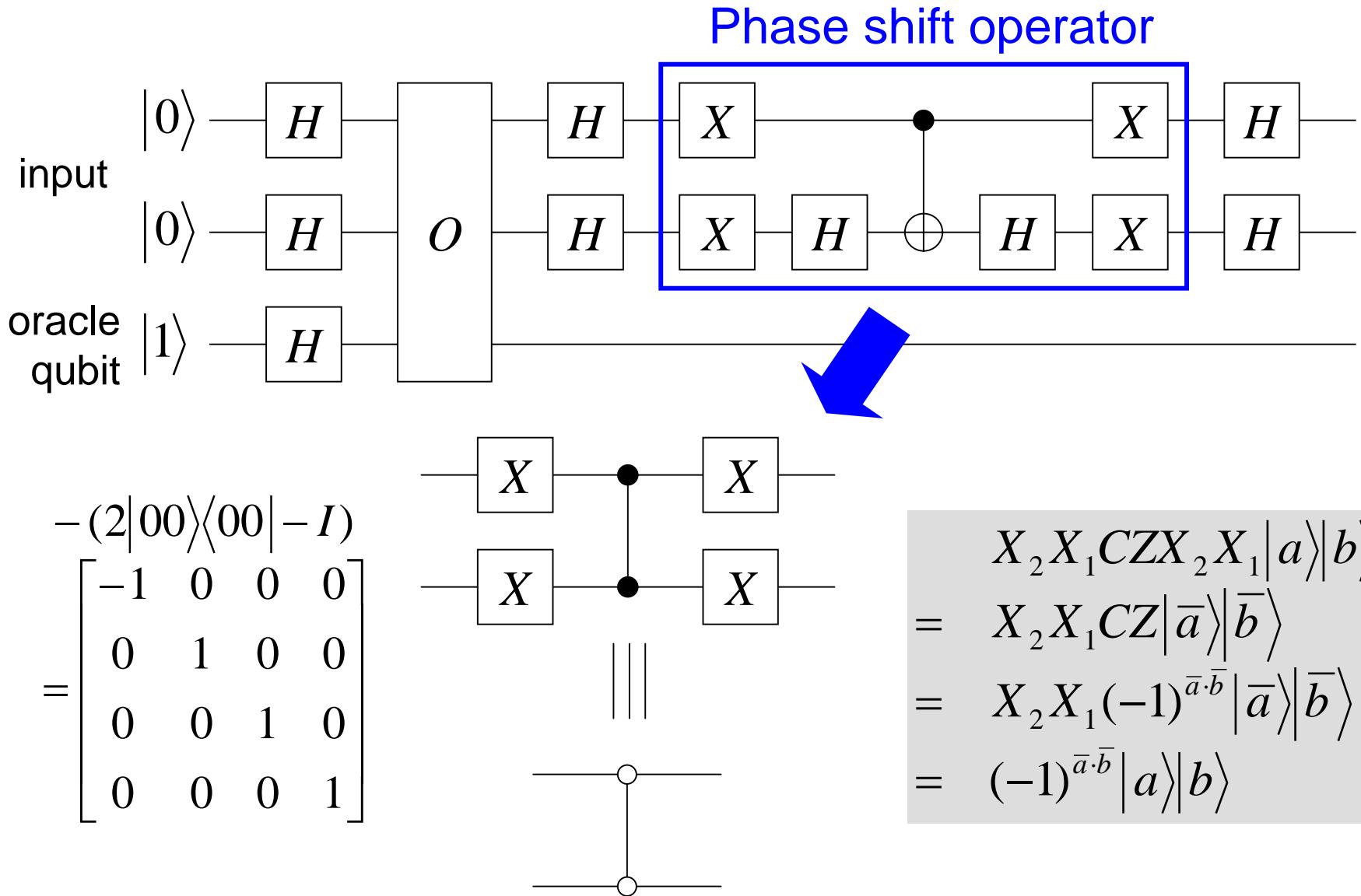
$$2\langle\alpha\rangle - \alpha_k = 2\frac{1}{4} \pm \frac{1}{2}$$

Oracle for 2-bit Grover

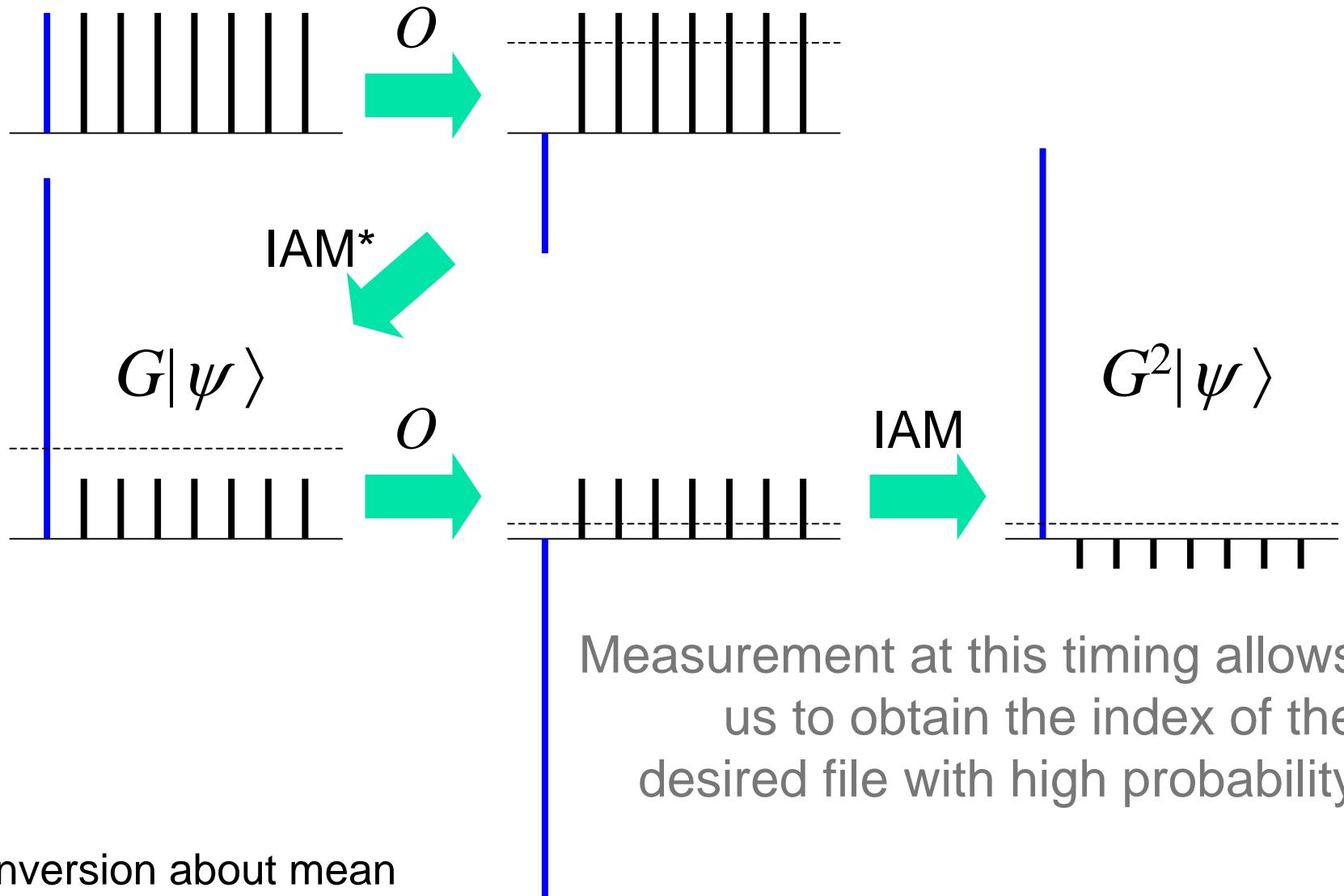


$$|x\rangle|q\rangle \xrightarrow{O} |x\rangle|q \oplus f(x)\rangle \quad f(x) = \begin{cases} 0 & (x \neq \beta) \\ 1 & (x = \beta) \end{cases}$$

Quantum circuit for 2-bit Grover



3-bit Grover



Geometric visualization

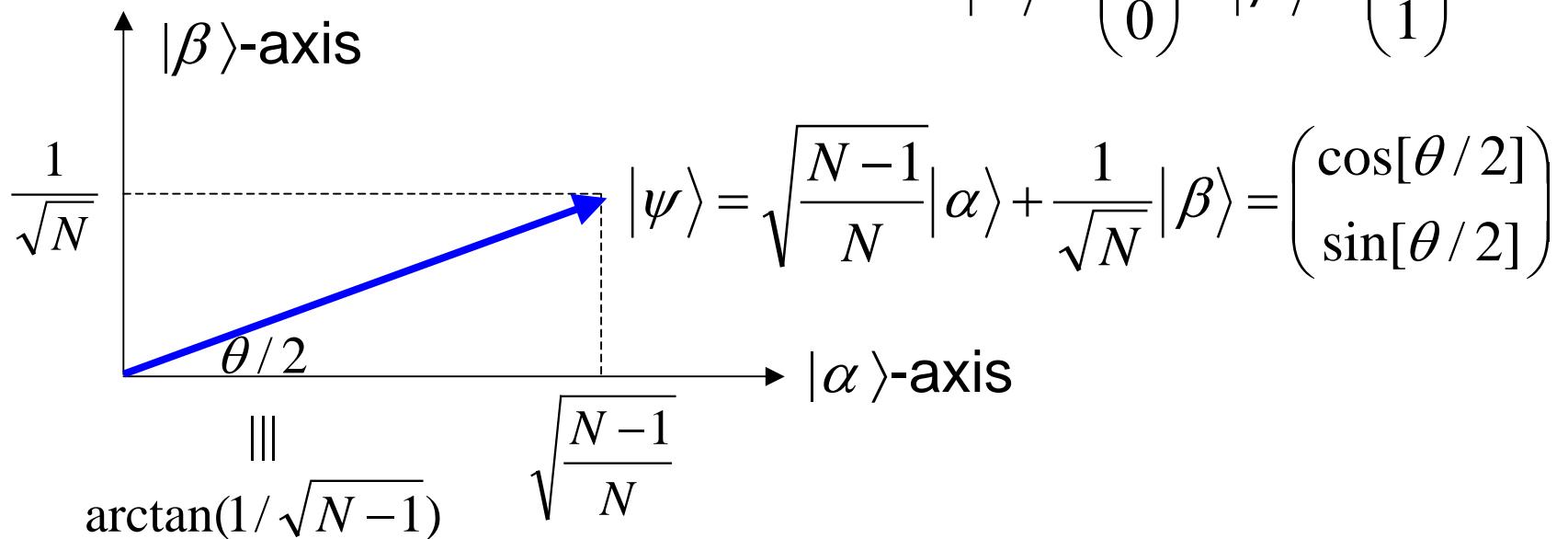
Sum over all x except β

$$|\alpha\rangle \equiv \frac{1}{\sqrt{N-1}} \sum_x' |x\rangle$$

$$\begin{aligned}\langle \alpha | \alpha \rangle &= \frac{1}{N-1} \sum_{x,x'}' \langle x' | x \rangle = 1 \\ \langle \alpha | \beta \rangle &= 0\end{aligned}$$

The initial state $|\psi\rangle$ is visualized as a vector in the real 2D plane spanned by $|\alpha\rangle$ and $|\beta\rangle$

$$|\alpha\rangle \equiv \begin{pmatrix} 1 \\ 0 \end{pmatrix} \quad |\beta\rangle \equiv \begin{pmatrix} 0 \\ 1 \end{pmatrix}$$

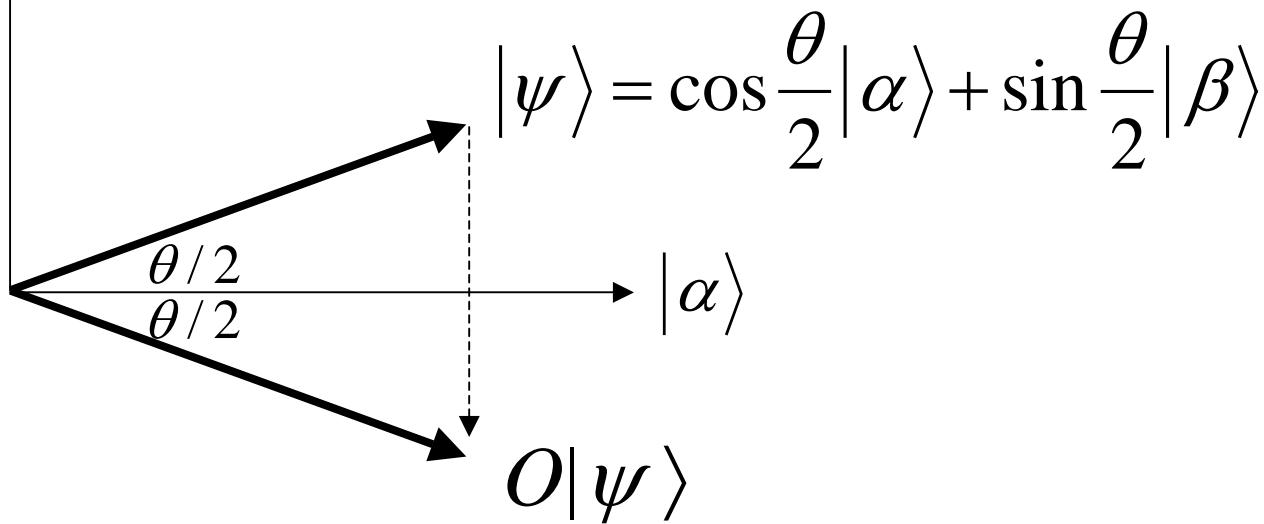


Geometric visualization

$$|\beta\rangle \quad \begin{cases} O|\alpha\rangle = |\alpha\rangle \\ O|\beta\rangle = -|\beta\rangle \end{cases} \Leftrightarrow O = \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix}$$

$$|x\rangle \xrightarrow{O} (-1)^{f(x)} |x\rangle$$
$$f(x) = \begin{cases} 0 & (x \neq \beta) \\ 1 & (x = \beta) \end{cases}$$

Reflection about the $|\alpha\rangle$ -axis



Geometric visualization

$$(2|\psi\rangle\langle\psi|-I)|\alpha\rangle = 2\cos\frac{\theta}{2}|\psi\rangle - |\alpha\rangle$$

$$|\psi\rangle = \cos\frac{\theta}{2}|\alpha\rangle + \sin\frac{\theta}{2}|\beta\rangle$$

$$\begin{aligned} &= \left(2\cos^2\frac{\theta}{2} - 1\right)|\alpha\rangle + 2\cos\frac{\theta}{2}\sin\frac{\theta}{2}|\beta\rangle \\ &= \cos\theta|\alpha\rangle + \sin\theta|\beta\rangle \end{aligned}$$

$$(2|\psi\rangle\langle\psi|-I)|\beta\rangle = 2\sin\frac{\theta}{2}|\psi\rangle - |\beta\rangle$$

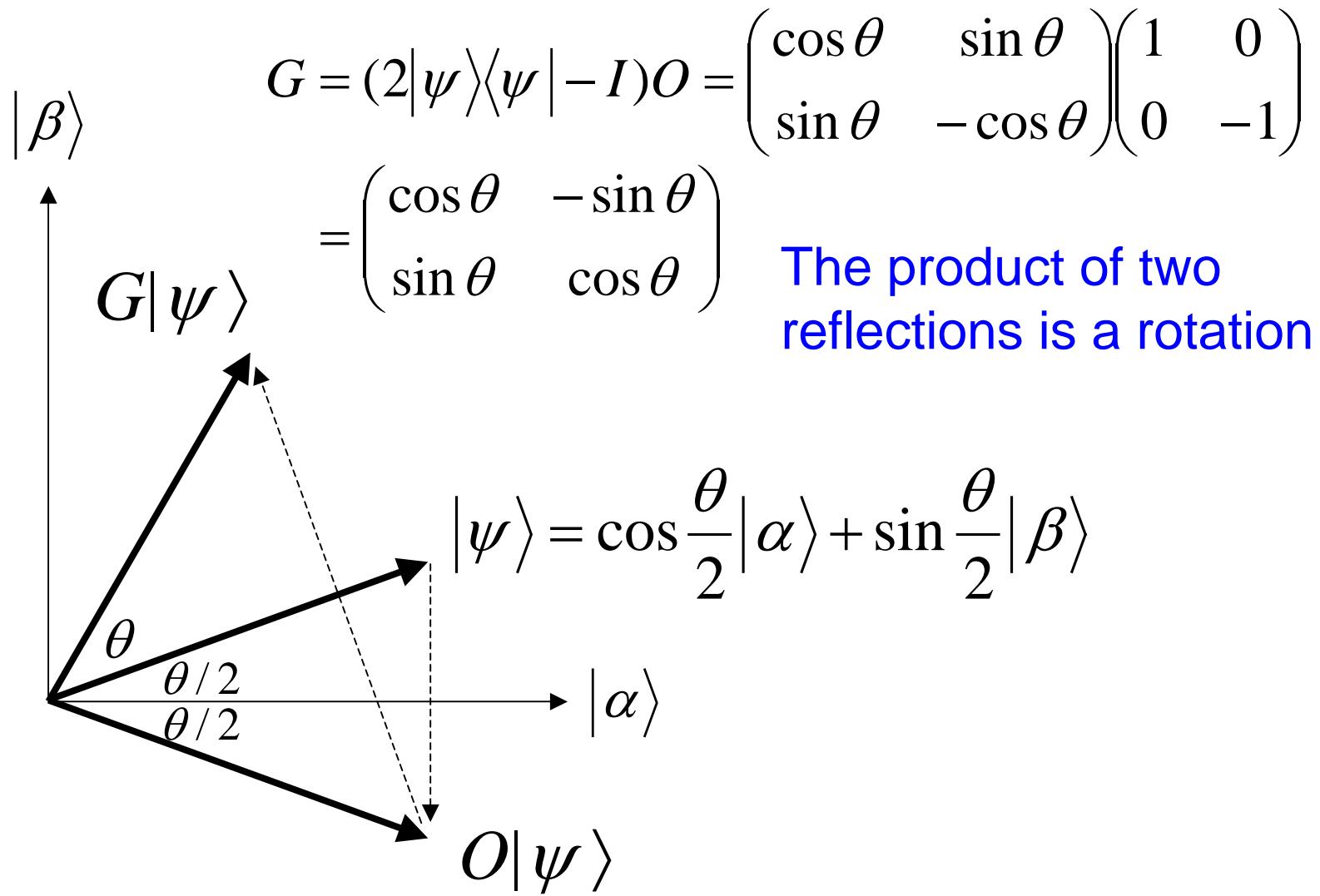
$$\begin{aligned} &= 2\cos\frac{\theta}{2}\sin\frac{\theta}{2}|\alpha\rangle + \left(2\sin^2\frac{\theta}{2} - 1\right)|\beta\rangle \\ &= \sin\theta|\alpha\rangle - \cos\theta|\beta\rangle \end{aligned}$$



$$2|\psi\rangle\langle\psi| - I = \begin{pmatrix} \cos\theta & \sin\theta \\ \sin\theta & -\cos\theta \end{pmatrix}$$

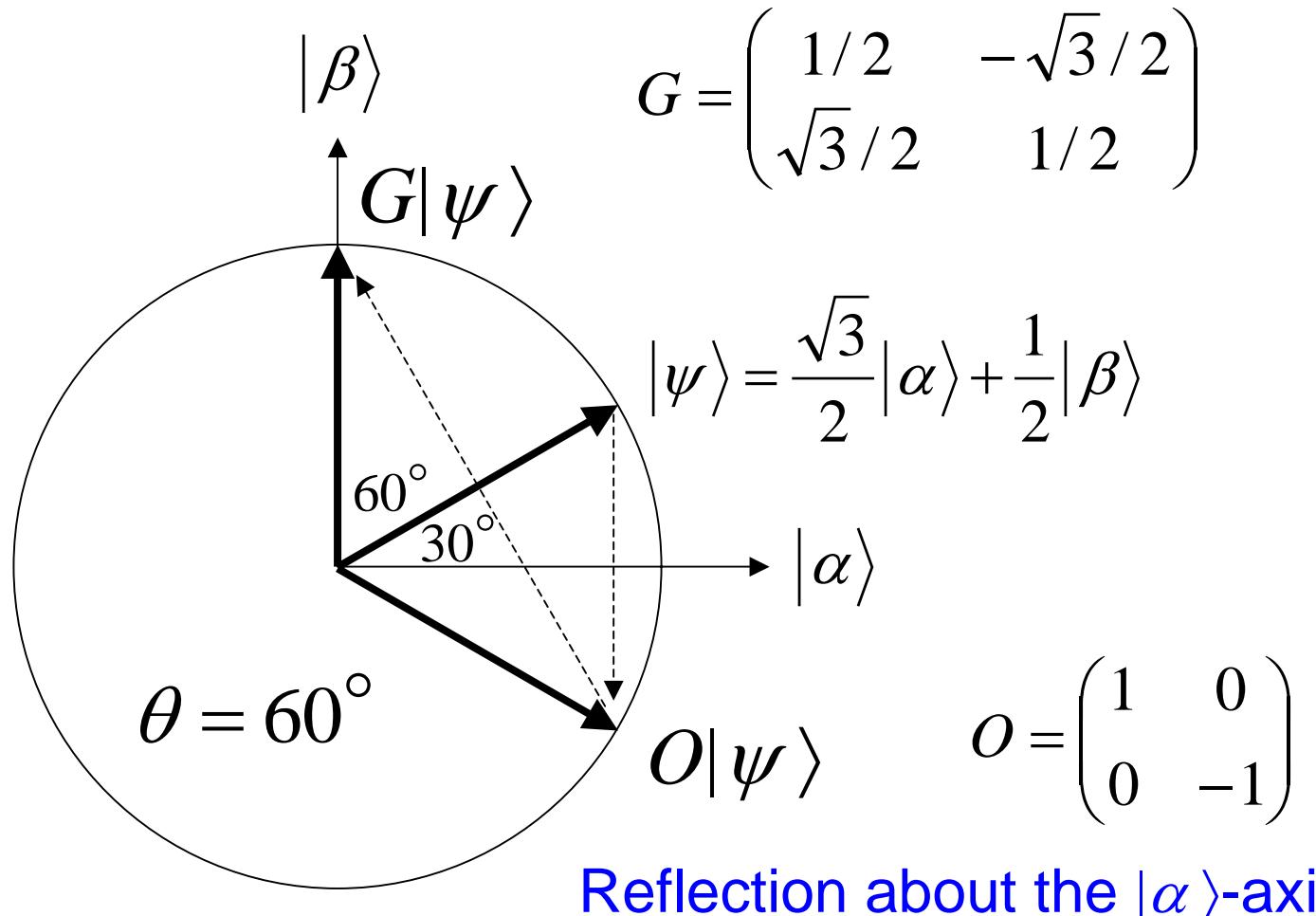
Reflection about
the vector $|\psi\rangle$

Geometric visualization

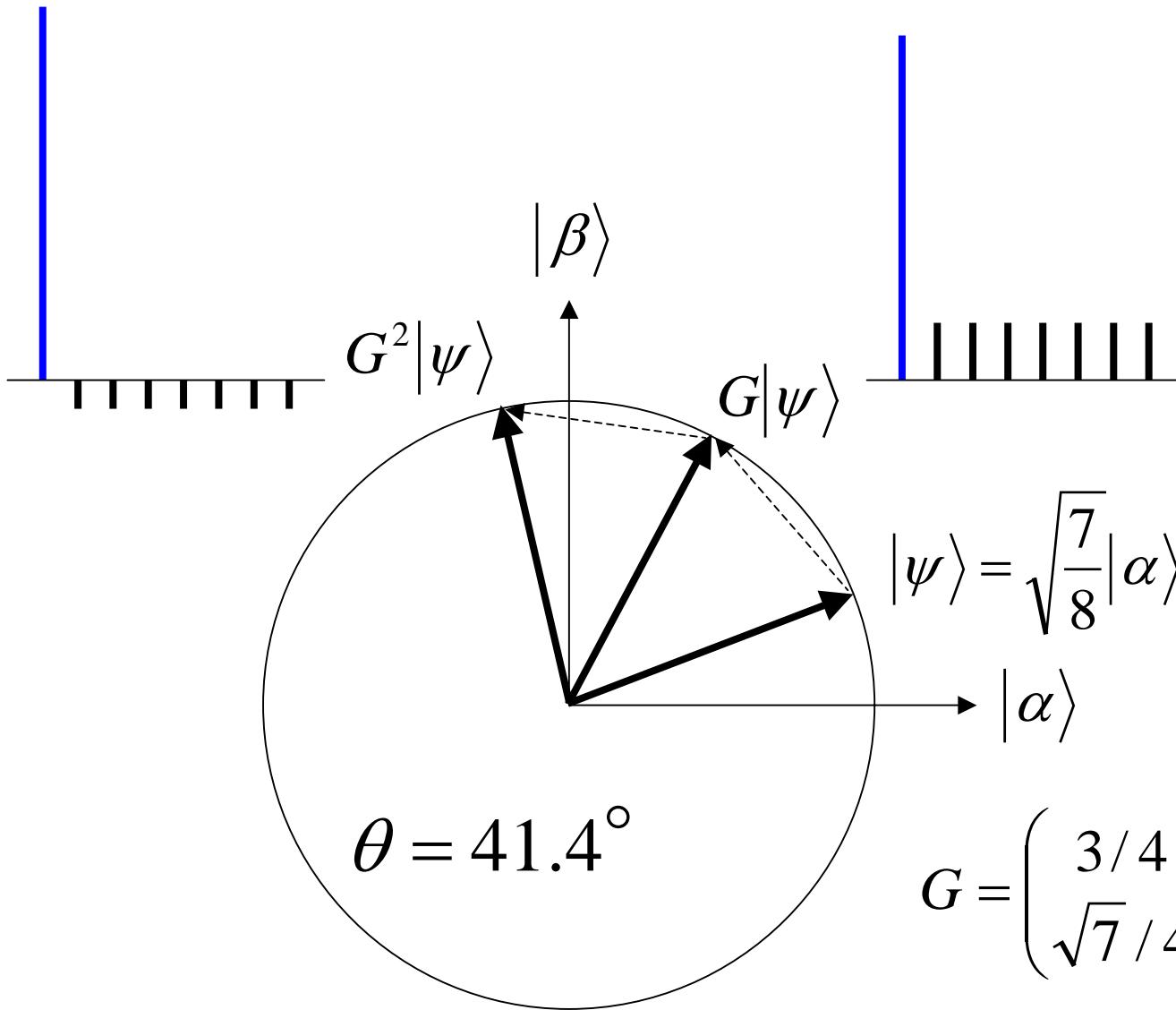


2-bit Grover

Rotation by 60 degrees



3-bit Grover



Performance

The state after repeating the Grover iteration k times

$$G^k |\psi\rangle = \cos\left(\frac{2k+1}{2}\theta\right) |\alpha\rangle + \sin\left(\frac{2k+1}{2}\theta\right) |\beta\rangle$$

We terminate the iteration when

$$\frac{2k'+1}{2}\theta \approx \frac{\pi}{2}$$

Assume θ is small (i.e., N large), then

$$\sin\frac{\theta}{2} = \frac{1}{\sqrt{N}} \approx \frac{\theta}{2}$$

The number of steps required to find the desired file

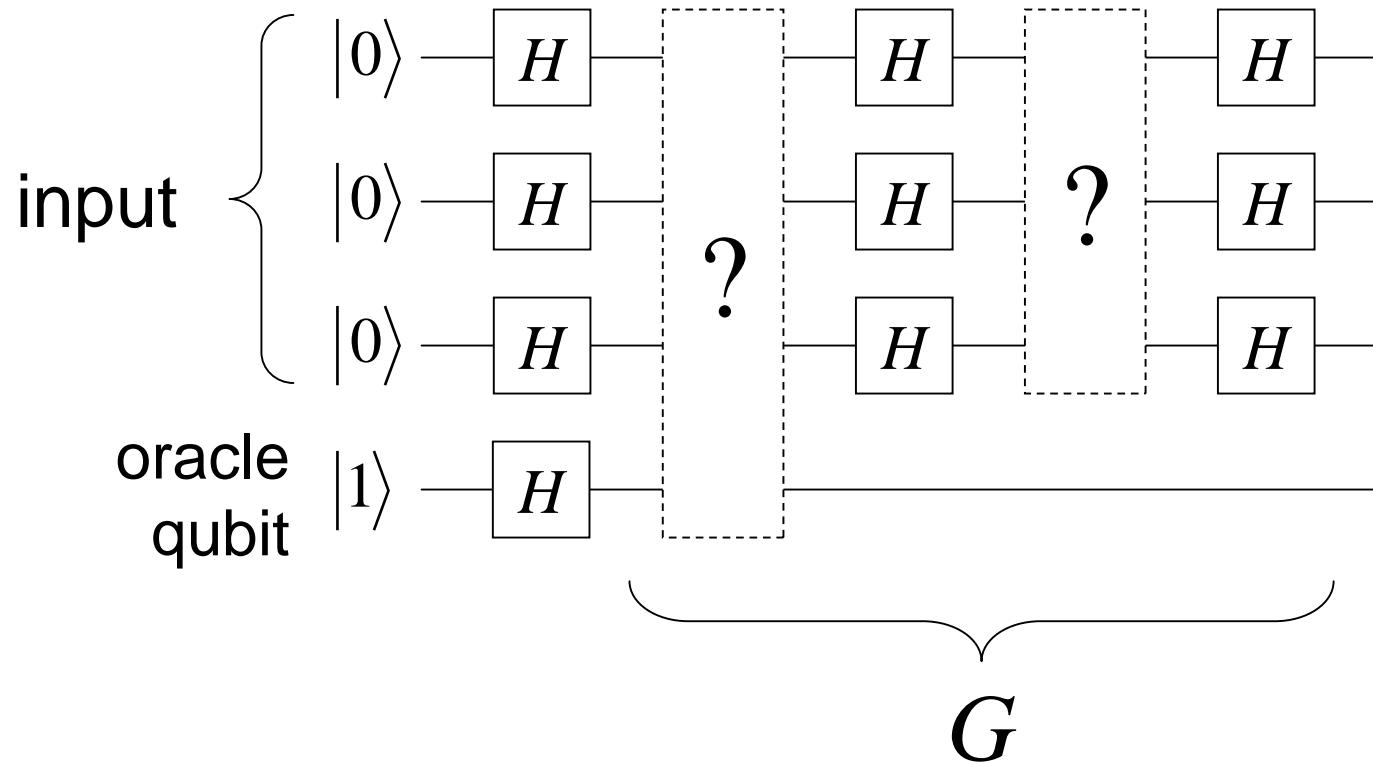
$$k' \approx \frac{\pi}{4} (\sqrt{N} - 1) = O(\sqrt{N})$$

Optimality

- Classical algorithms take $O(N)$ operations for searching N items
- Grover's algorithm can search N items by calling the oracle only $O(N^{1/2})$ times
- It is shown that Grover's algorithm is optimal, i.e., any quantum algorithms require at least $O(N^{1/2})$ times oracle callings for searching
- The proof is beyond the scope of this introductory lecture

Quiz

Complete the quantum circuit for 3-bit Grover to search the file #7



Answer

